

BACKEND DESIGN

Placement and Pin Assignment

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Placement

Problem Definition

- **Input:**
 - A set of blocks, both fixed and flexible.
 - Area of the block $A_i = w_i \times h_i$
 - Constraint on the shape of the block (rigid/flexible)
 - Pin locations of fixed blocks.
 - A netlist.
- **Requirements:**
 - Find locations for each block so that no two blocks overlap.
 - Determine shapes of flexible blocks.
- **Objectives:**
 - Minimize area.
 - Reduce wire-length for critical nets.

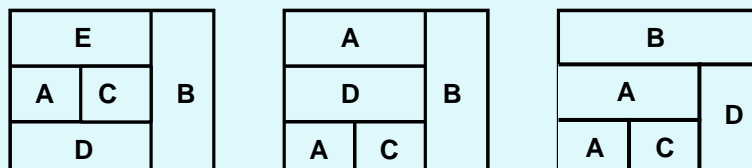
Difference Between Floorplanning and Placement

- The problems are similar in nature.
- **Main differences:**
 - In floorplanning, some of the blocks may be flexible, and the exact locations of the pins not yet fixed.
 - In placement, all blocks are assumed to be of well-defined geometrical shapes, with defined pin locations.
- **Points to note:**
 - Floorplanning problem is more difficult as compared to placement.
 - Multiple choice for the shape of a block.
 - In some of the VLSI design styles, the two problems are identical.

An Example for Rigid Blocks

Module	Width	Height
A	1	1
B	1	3
C	1	1
D	1	2
E	2	1

Some of the Feasible Floorplans



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Design Style Specific Issues

- **Full Custom**
 - All the steps required for general cells.
- **Standard Cell**
 - Dimensions of all cells are fixed.
 - Floorplanning problem is simply the placement problem.
 - For large netlists, two steps:
 - First do global partitioning.
 - Placement for individual regions next.
- **Gate Array**
 - Floorplanning problem same as placement problem.

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Estimating Cost of a Floorplan

- The number of feasible solutions of a floorplanning problem is very large.
 - Finding the best solution is NP-hard.
- Several criteria used to measure the quality of floorplans:
 - a) Minimize area
 - b) Minimize total length of wire
 - c) Maximize routability
 - d) Minimize delays
 - e) Any combination of above

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Contd.

- How to determine area?
 - Not difficult.
 - Can be easily estimated because the dimensions of each block is known.
 - Area A computed for each candidate floorplan.
- How to determine wire length?
 - A coarse measure is used.
 - Based on a model where all I/O pins of the blocks are merged and assumed to reside at its center.
 - Overall wiring length $L = \sum_{i,j} (c_{ij} * d_{ij})$
 - where c_{ij} : connectivity between blocks i and j
 - d_{ij} : Manhattan distances between the centres of rectangles of blocks i and j

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Contd.

- **Typical cost function used:**

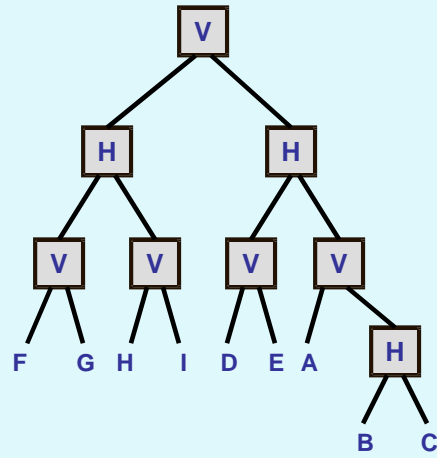
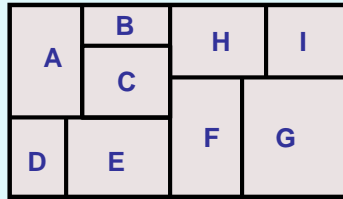
$$\text{Cost} = w1 * A + w2 * L$$

where $w1$ and $w2$ are user-specified parameters.

Slicing Structure

- **Definition**
 - A rectangular dissection that can be obtained by repeatedly splitting rectangles by horizontal and vertical lines into smaller rectangles.
- **Slicing Tree**
 - A binary tree that models a slicing structure.
 - Each node represents a *vertical cut line (V)*, or a *horizontal cut line (H)*.
 - A third kind of node called *Wheel (W)* appears for non-sliceable floorplans (discussed later).
 - Each leaf is a basic block (rectangle).

A Slicing Floorplan and its Slicing Tree

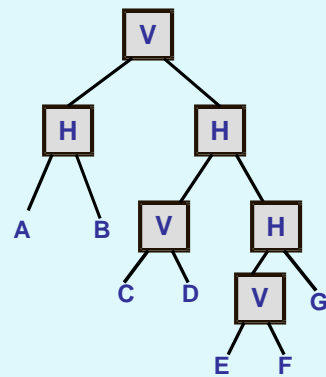
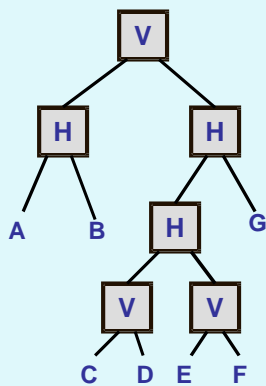
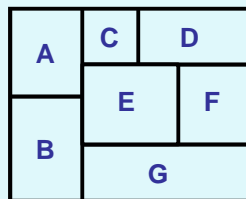


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Slicing Tree is not Unique

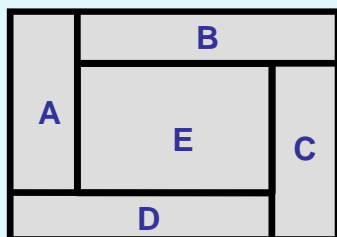


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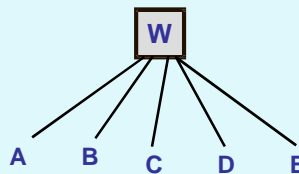
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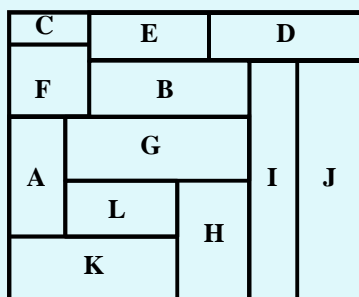
A Non-Slicing Floorplan



Also called "WHEEL"



A Hierarchical Floorplan



The dissection tree will contain "wheel".

Floorplanning Algorithms

- Several broad classes of algorithms:
 - Integer programming based
 - Rectangular dual graph based
 - Hierarchical tree based
 - Simulated annealing based
 - Other variations

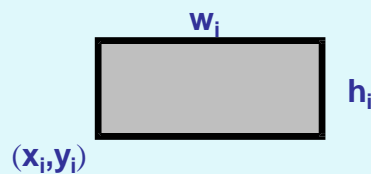
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Integer Linear Programming Formulation

- The problem is modeled as a set of linear equations using 0/1 integer variables.
- Given:
 - Set of n blocks $S = \{B_1, B_2, \dots, B_n\}$ which are rigid and have fixed orientation.
 - 4-tuple associated with each block
 (x_i, y_i, w_i, h_i)



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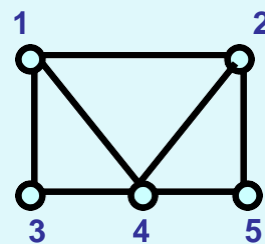
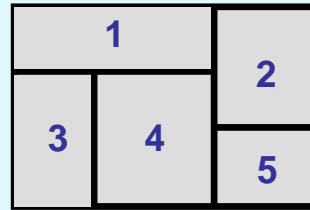
Mathematical Formulation

Rectangular Dual-Graph Approach

- **Basic Concept:**
 - Output of partitioning algorithms represented by a graph.
 - Floorplans can be obtained by converting the graph into its rectangular dual.
- **The rectangular dual of a graph satisfies the following properties:**
 - Each vertex corresponds to a distinct rectangle.
 - For every edge, the corresponding rectangles are adjacent.

A Rectangular Floorplan & its Dual Graph

- Without loss of generality, we assume that a rectangular floorplan contains no cross junctions.
- Under this assumption, the dual graph of a rectangular floorplan is a planar triangulated graph (PTG).



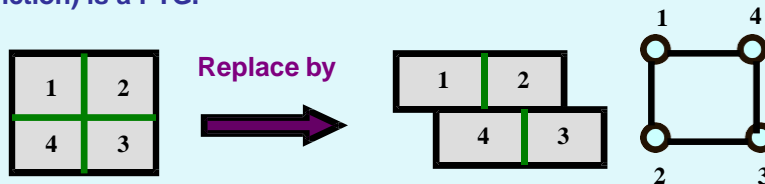
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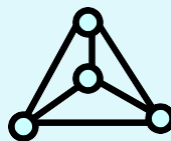
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Contd.

- Every dual graph of a rectangular floorplan (without cross junction) is a PTG.



- However, not every PTG corresponds to a rectangular floorplan.



Complex triangle

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Drawbacks

- A new approach to floorplanning, in which many sub-problems are still unsolved.
- The main problem concerns the existence of the rectangular dual, i.e. the elimination of complex triangles.
 - Select a minimum set E of edges such that each complex triangle has at least one edge in E.
 - A vertex can be added to each edge of E to eliminate all complex triangles.
 - The weighted complex triangle elimination problem has been shown to be NP-complete.
 - Some heuristics are available.

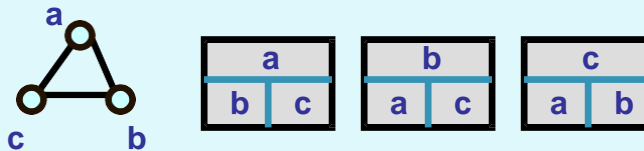
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Hierarchical Approach

- Widely used approach to floorplanning.
 - Based on a divide-and-conquer paradigm.
 - At each level of the hierarchy, only a small number of rectangles are considered.
- A small graph, and all possible floorplans.



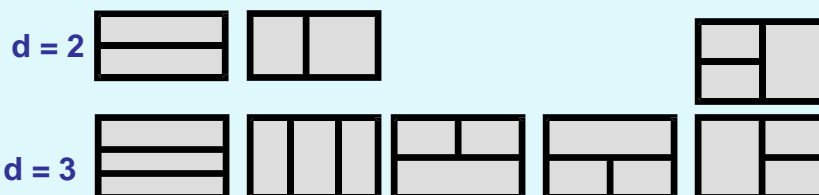
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- After an optimal configuration for the three modules has been determined, they are merged into a larger module.
 - The vertices 'a', 'b', 'c' are merged into a super vertex at the next level.

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- The number of floorplans increases exponentially with the number of modules 'd' considered at each level.
 - 'd' is thus limited to a small number (typically $d < 6$).
 - All possible floorplans for:



Hierarchical Approach :: Bottom-Up

- Hierarchical approach works best in bottom-up fashion.
- Modules are represented as vertices of a graph, while edges represent connectivity.
 - Modules with high connectivity are clustered together.
 - Number of modules in each cluster $\leq d$.
 - An optimal floorplan for each cluster is determined by exhaustive enumeration.
 - The cluster is merged into a larger module for high-level processing.

Contd.

- A Greedy Procedure
 - Sort the edges in decreasing weights.
 - The heaviest edge is chosen, and the two modules of the edge are clustered in a greedy fashion.
 - Restriction: number of modules in each cluster $\leq d$.
 - In the next higher level, vertices in a cluster are merged, and edge weights are summed up accordingly.

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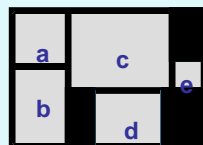
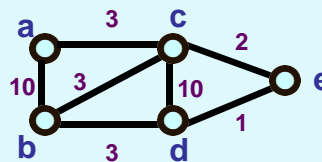
- **Problem**
 - Some lightweight edges may be chosen at higher levels in the hierarchy, resulting in adjacency of two clusters of highly incompatible areas.
- **Possible solution**
 - Arbitrarily assign a small cluster to a neighboring cluster when their sizes will be too small for processing at a higher level of the hierarchy.

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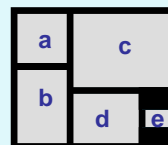
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Example



Greedy clustering



Merging small clusters

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Hierarchical Approach :: Top-Down

- The fundamental step is the partitioning of modules.
 - Each partition is assigned to a child floorplan.
 - Partitioning is recursively applied to the child floorplans.
- Major issue here is to obtain balanced graph partitioning.
 - k-way partitioning, in general.
- Not very widely used due to the difficulty of obtaining balanced partitions.

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- One can combine top-down and bottom-up approaches.
 - Apply bottom-up technique to obtain a set of clusters.
 - Apply top-down approach to these clusters.

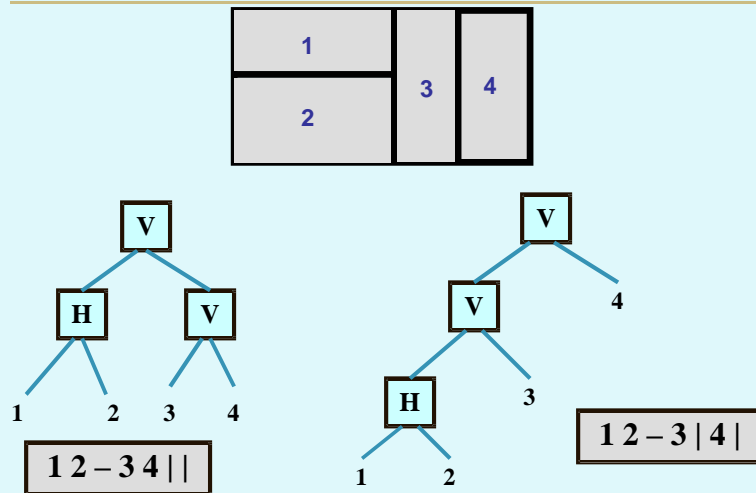
Simulated Annealing

- **Important issues in the design of a simulated annealing optimization problem:**
 1. The solution space.
 2. The movement from one solution to another.
 3. The cost evaluation function.
- **A solution by Wong & Liu is applicable to sliceable floorplans only.**
 - Floorplan can be represented by a tree.
 - Postfix notations used for easy representation and manipulation.

Some Notations

- **Dissection operators defined:**
 - ijH means rectangle j is on top of rectangle i.
 - ijV means rectangle j is on left of rectangle i.
- **Normalized Polish expression:**
 - A Polish expression corresponding to a floorplan is called normalized if it has no consecutive H's or V's.
 - Used for the purpose of removing redundant solutions from the solution space.
 - There may be several several Polish expressions that correspond to the same slicing floorplan.
 - We have seen earlier that the same floorplan can have more than one slicing tree.

Example: Two tree representations of a floorplan



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Simulated Annealing

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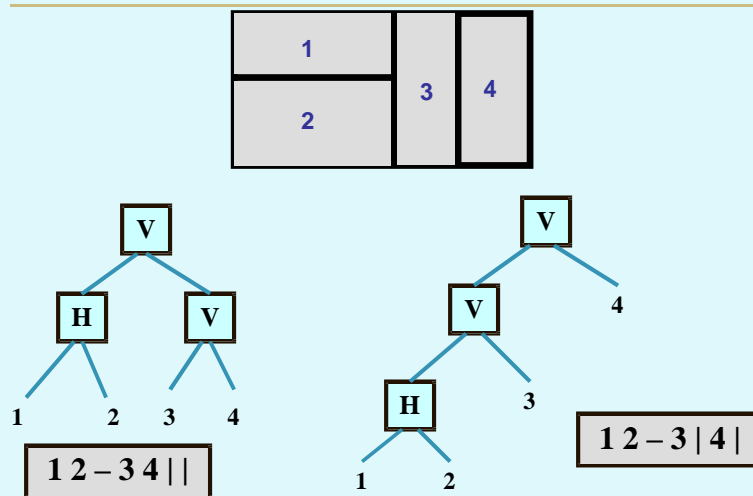
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Example: Two tree representations of a floorplan



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Parameters of the Algorithm

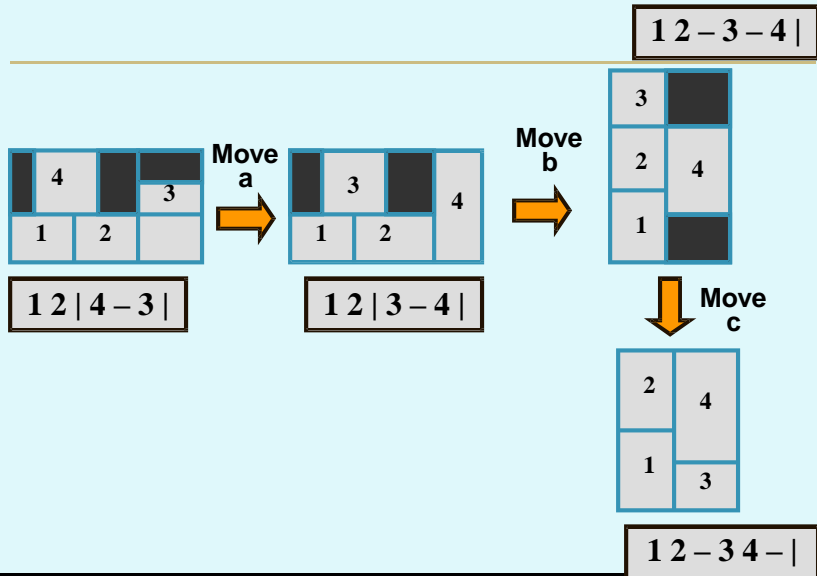
- **Solution Perturbations (Move)**
 - a) Swap two adjacent operands
 - b) Complement a series of operators between two operands (called a *chain*).
 $V=H$ and $H=V$
 - c) Swap two adjacent operand and operator.
- **We accept a move only it results in a normalized expression.**
 - Only move “c” may result in a non-normalized solution.
 - We need to check only when move “c” is applied.

Contd.

- **Cost Function:**
 - Typical cost function used:
$$\text{Cost} = w_1 * A + w_2 * L$$

where A is the area of the smallest rectangle enveloping the given basic rectangles,
L is the overall interconnection length,
 w_1 and w_2 are user-specified parameters.

An Example



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Pin Assignment

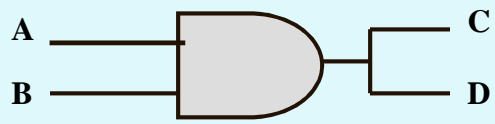
Introduction

- The purpose is to define the signal that each pin will receive.
- It can be done
 - During floorplanning
 - During placement
 - After placement is fixed
- For undesigned blocks, a good assignment of pins improves placement.
- If the blocks are already designed, still some pins can be exchanged.

Pin Assignment

- Input:
 - A placement of blocks.
 - Number of pins on each block, possibly an ordering.
 - A netlist.
- Requirements:
 - To determine the pin locations on the blocks.
- Objectives:
 - To minimize net-length.

-
- **Functionally equivalent pins:**
 - Exchanging the signals does not affect the circuit.
 - **Equipotential pins:**
 - Both are internally connected and represent the same net.

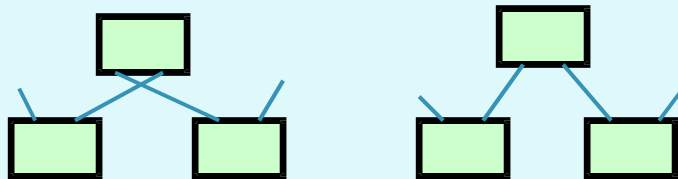


A,B :: functionally equivalent

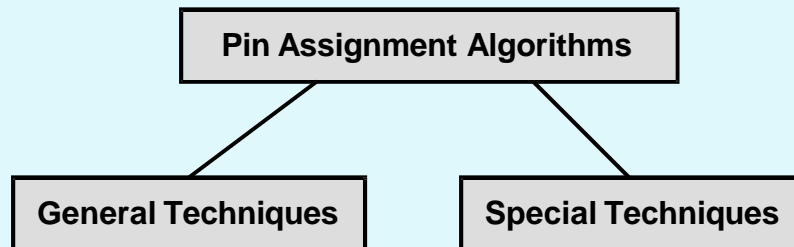
C,D :: equipotential

Problem Formulation

- Purpose is to optimize the assignment of nets within a functionally equivalent (or equipotential) pin groups.
- **Objective:**
 - To reduce congestion or reduce the number of crossovers.



Classification of Algorithms



Concentric Circle Mapping

Topological Method

Nine Zone Method

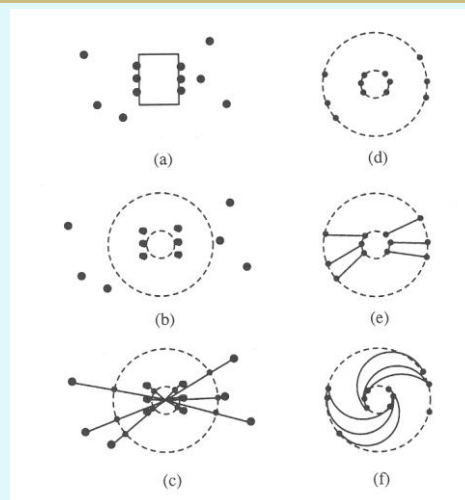
Channel Pin Assignment

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Concentric Circle Mapping



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Topological Pin Assignment

- Similar to concentric circle mapping.
- Easier to complete pin assignment.
 - When there is interference from other components and barriers.
 - For nets connected to more than two pins.
- If a net has been assigned to more than two pins, then the pin closest to the center of the primary component is chosen.
- Pins of primary component are mapped onto a circle as before.
- Beginning at the bottom of the circle, and moving clockwise, the pins are assigned to nets.

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Nine Zone Method

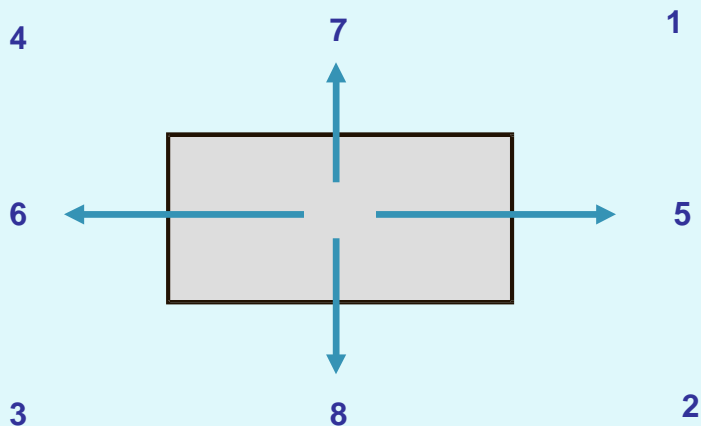
- Based on zones in a Cartesian co-ordinate system.
- The center of the co-ordinate system is located inside a group of interchangeable pins on a component.
 - This component is called *pin class*.
- A net rectangle is defined by each of the nets connected to the pin class.
 - There are nine zones in which this rectangle can be positioned.

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The Nine Pin Zones



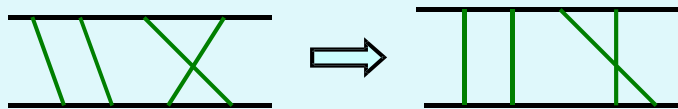
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Channel Pin Assignment

- A significant portion of the chip area is used for channel routing.
 - After the placement phase, the position of terminals on the boundaries of a block are not fixed.
 - They may be moved before routing begins.



- Yang & Wong proposed a dynamic programming formulation to the channel pin assignment problem.

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Integrated Approach

- **Better understanding of the different stages in physical design automation over the years.**
 - Attempts are being made to merge some steps of the design cycle.
 - For example, floorplanning and placement are considered together.
 - Sometimes, placement and routing stages can also be combined together.
- **Still a problem of research.**