

# **Lecture 10: Circuit Families**

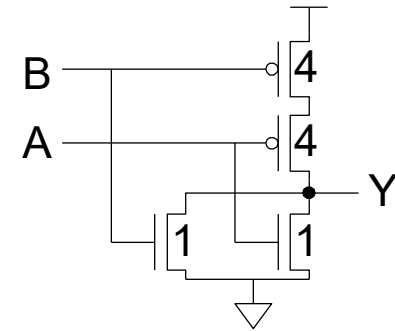
# Outline

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- ☐ Pseudo-nMOS Logic
- ☐ Dynamic Logic
- ☐ Pass Transistor Logic

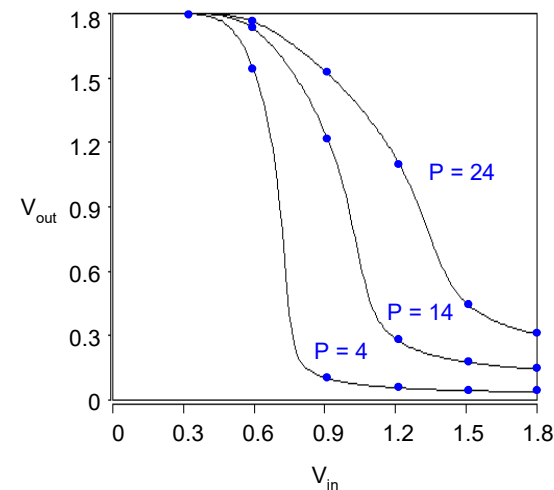
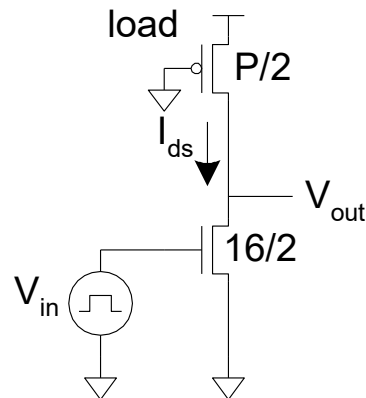
# Introduction

- ❑ What makes a circuit fast?
  - $I = C \, dV/dt \rightarrow t_{pd} \propto (C/I) \Delta V$
  - low capacitance
  - high current
  - small swing
- ❑ Logical effort is proportional to  $C/I$
- ❑ pMOS are the enemy!
  - High capacitance for a given current
- ❑ Can we take the pMOS capacitance off the input?
- ❑ Various circuit families try to do this...



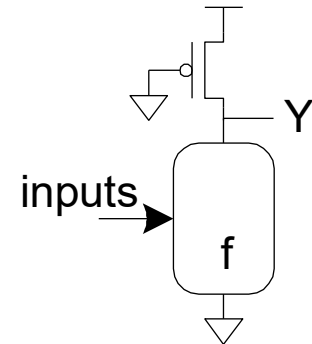
# Pseudo-nMOS

- ❑ In the old days, nMOS processes had no pMOS
  - Instead, use pull-up transistor that is always ON
- ❑ In CMOS, use a pMOS that is always ON
  - *Ratio* issue
  - Make pMOS about  $\frac{1}{4}$  effective strength of pulldown network

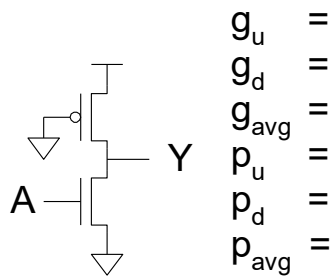


# Pseudo-nMOS Gates

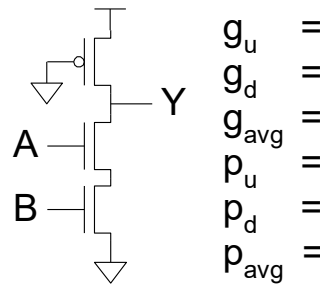
- ❑ Design for unit current on output to compare with unit inverter.
- ❑ pMOS fights nMOS



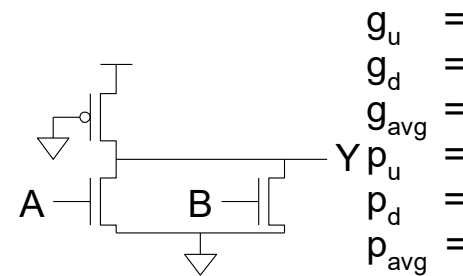
Inverter



NAND2

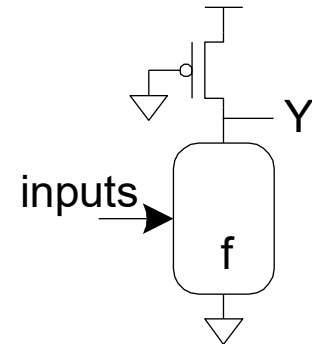


NOR2

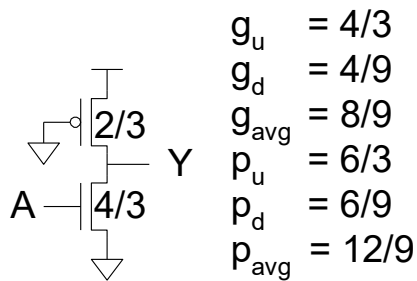


# Pseudo-nMOS Gates

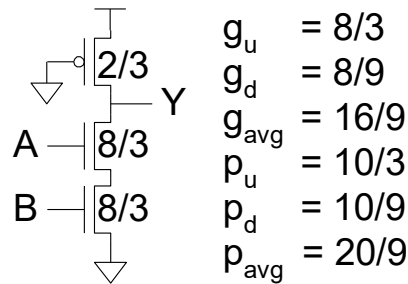
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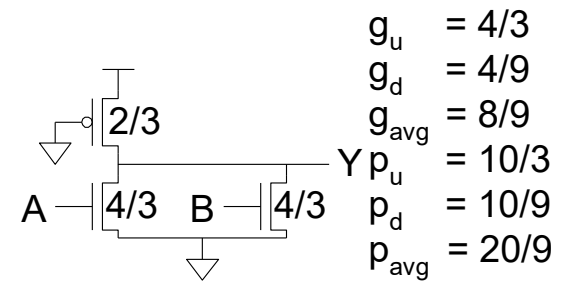
Inverter



NAND2



NOR2



# Pseudo-nMOS Design

- ❑ Ex: Design a k-input AND gate using pseudo-nMOS. Estimate the delay driving a fanout of H

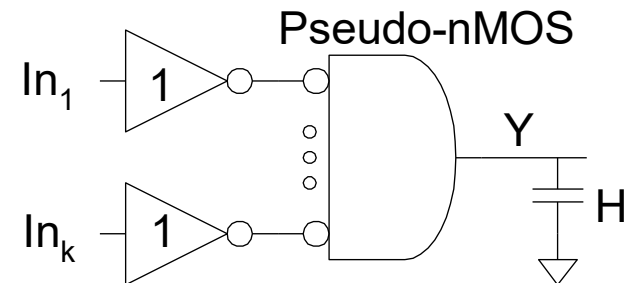
- ❑  $G =$

- ❑  $F =$

- ❑  $P =$

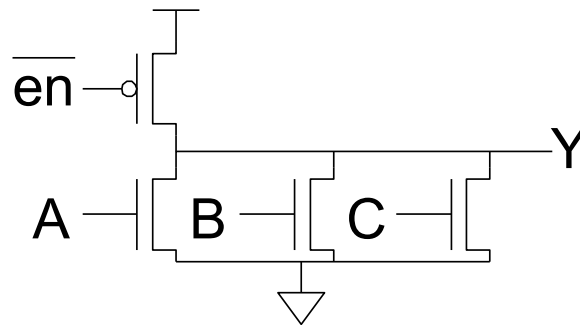
- ❑  $N =$

- ❑  $D =$



# Pseudo-nMOS Power

- ❑ Pseudo-nMOS draws power whenever  $Y = 0$ 
  - Called static power  $P = I_{DD} V_{DD}$
  - A few mA / gate \* 1M gates would be a problem
  - Explains why nMOS went extinct
- ❑ Use pseudo-nMOS sparingly for wide NORs
- ❑ Turn off pMOS when not in use





# Ratio Example

- ❑ The chip contains a 32 word x 48 bit ROM
  - Uses pseudo-nMOS decoder and bitline pullups
  - On average, one wordline and 24 bitlines are high
- ❑ Find static power drawn by the ROM
  - $I_{\text{on-p}} = 36 \mu\text{A}$ ,  $V_{\text{DD}} = 1.0 \text{ V}$

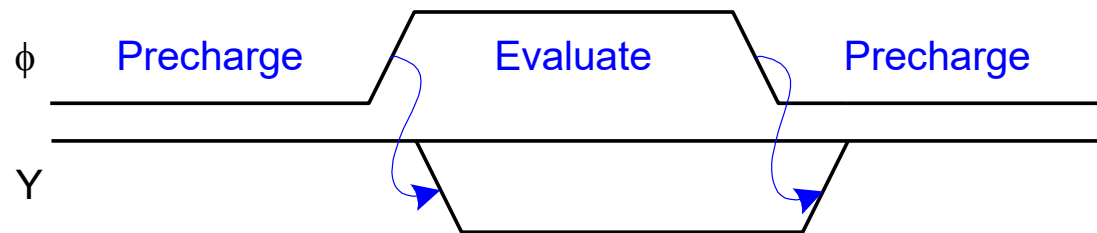
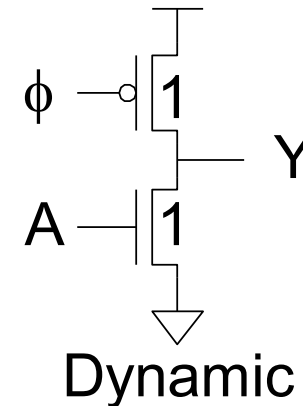
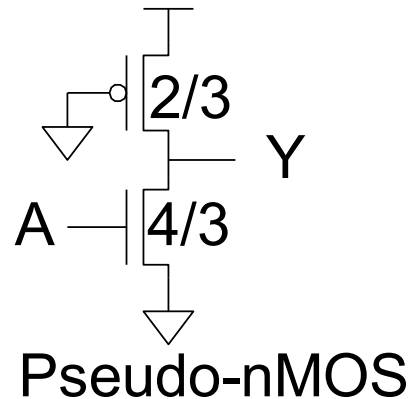
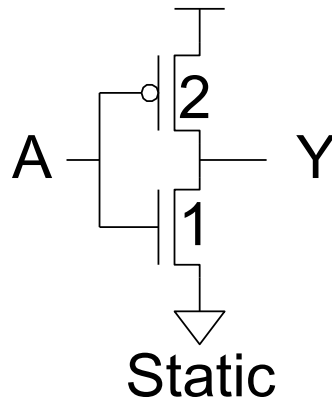
❑ Solution:

$$P_{\text{pull-up}} =$$

$$P_{\text{static}} =$$

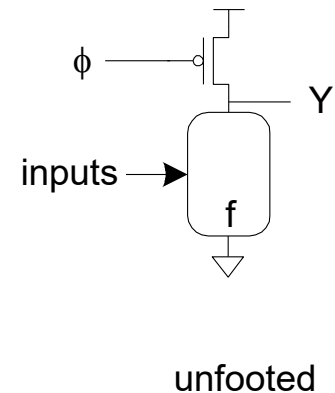
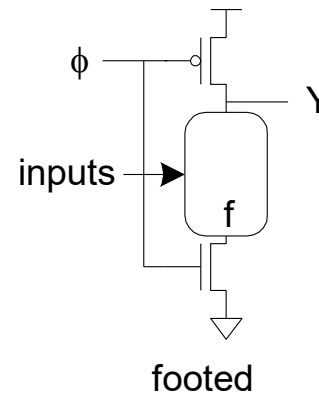
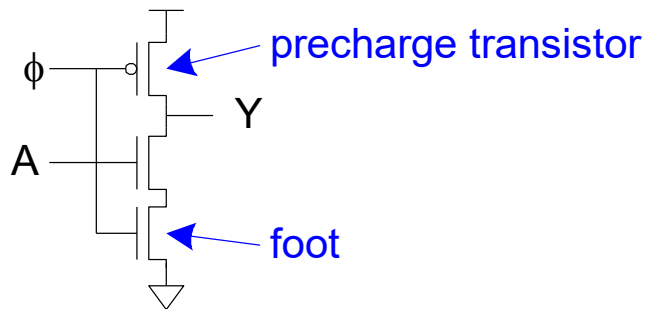
# Dynamic Logic

- ❑ *Dynamic* gates uses a clocked pMOS pullup
- ❑ Two modes: *precharge* and *evaluate*



# The Foot

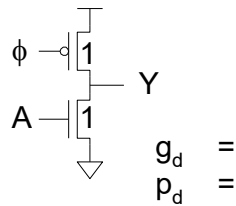
- ❑ What if pulldown network is ON during precharge?
- ❑ Use series evaluation transistor to prevent fight.



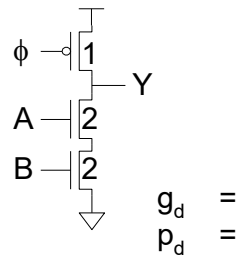
# Logical Effort

unfooted

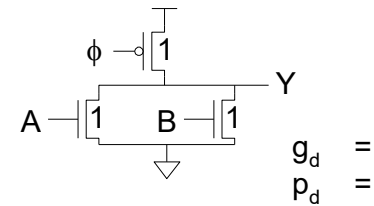
Inverter



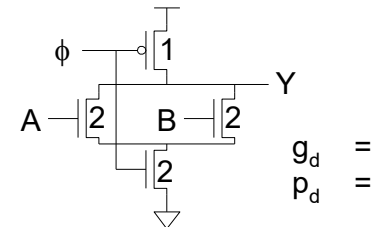
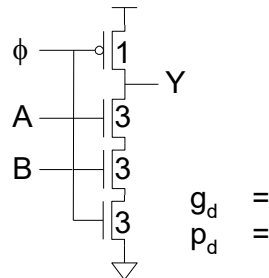
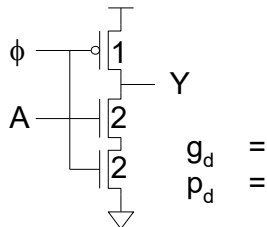
NAND2



NOR2



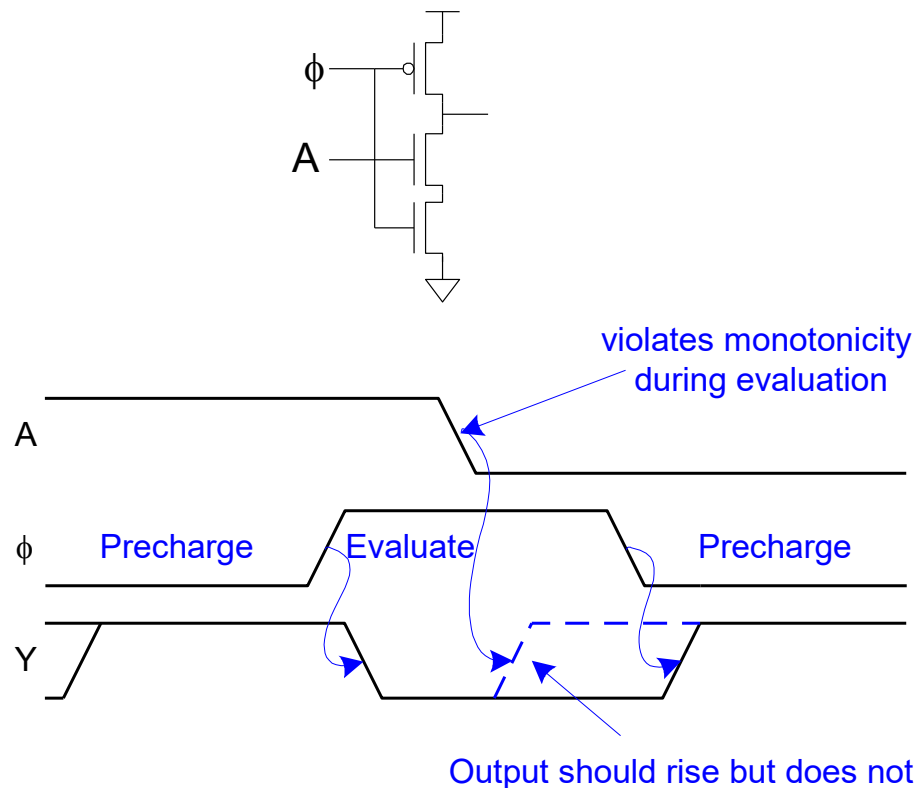
footed



# Monotonicity

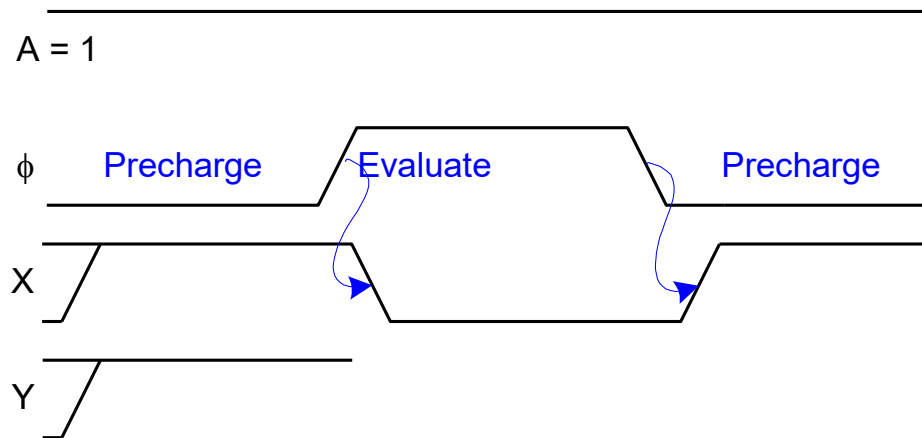
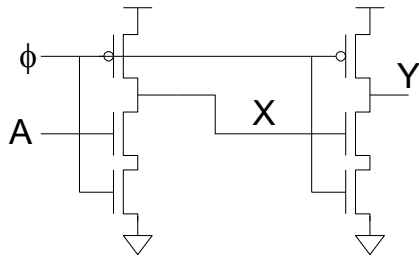
- ❑ Dynamic gates require *monotonically rising* inputs during evaluation

- 0  $\rightarrow$  0
- 0  $\rightarrow$  1
- 1  $\rightarrow$  1
- But not 1  $\rightarrow$  0



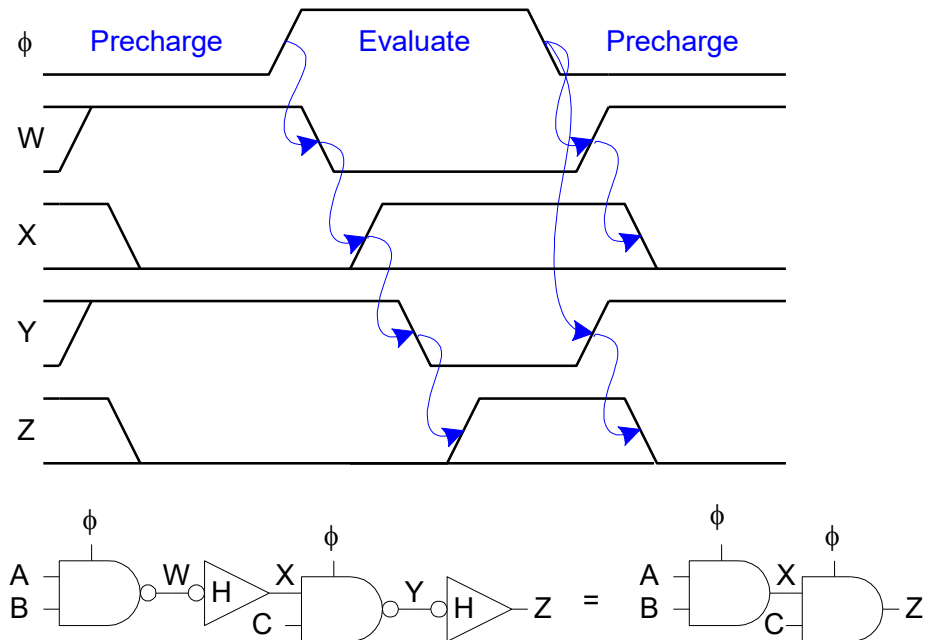
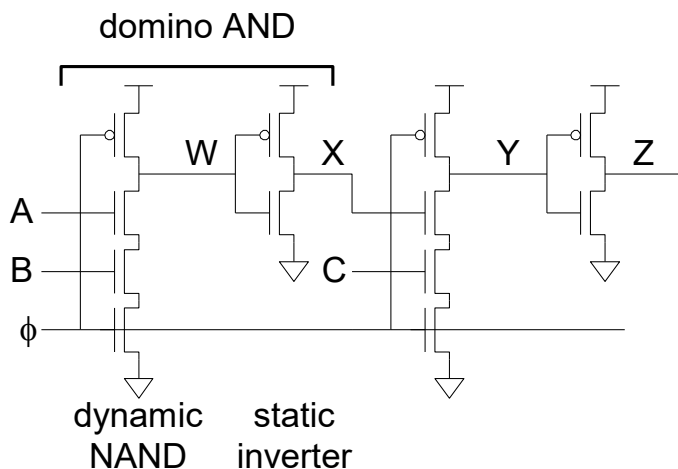
# Monotonicity Woes

- ❑ But dynamic gates produce monotonically falling outputs during evaluation
- ❑ Illegal for one dynamic gate to drive another!



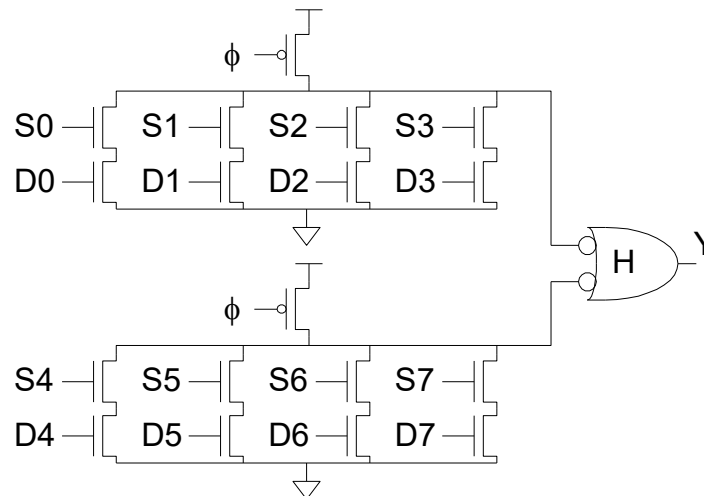
# Domino Gates

- ❑ Follow dynamic stage with inverting static gate
  - Dynamic / static pair is called domino gate
  - Produces monotonic outputs



# Domino Optimizations

- ❑ Each domino gate triggers next one, like a string of dominos toppling over
- ❑ Gates evaluate sequentially but precharge in parallel
- ❑ Thus evaluation is more critical than precharge
- ❑ HI-skewed static stages can perform logic

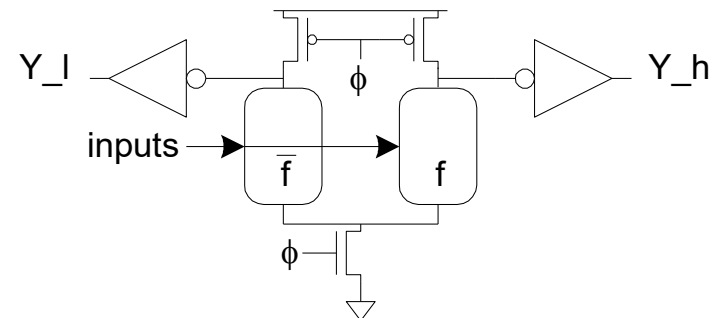




# Dual-Rail Domino

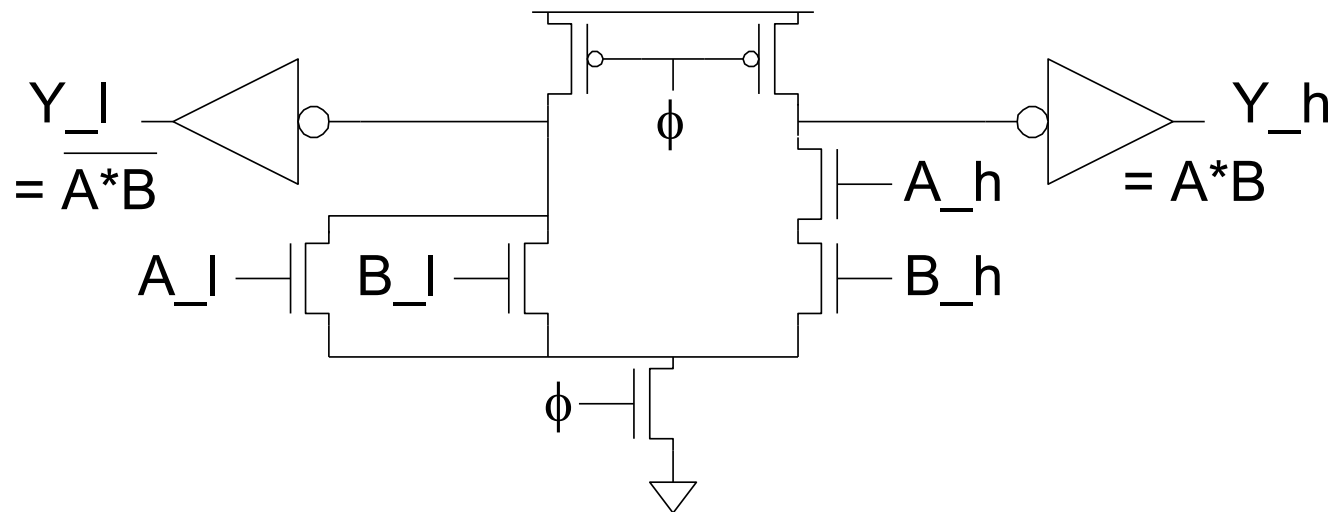
- ❑ Domino only performs noninverting functions:
  - AND, OR but not NAND, NOR, or XOR
- ❑ Dual-rail domino solves this problem
  - Takes true and complementary inputs
  - Produces true and complementary outputs

sig_h	sig_l	Meaning
0	0	Precharged
0	1	'0'
1	0	'1'
1	1	invalid



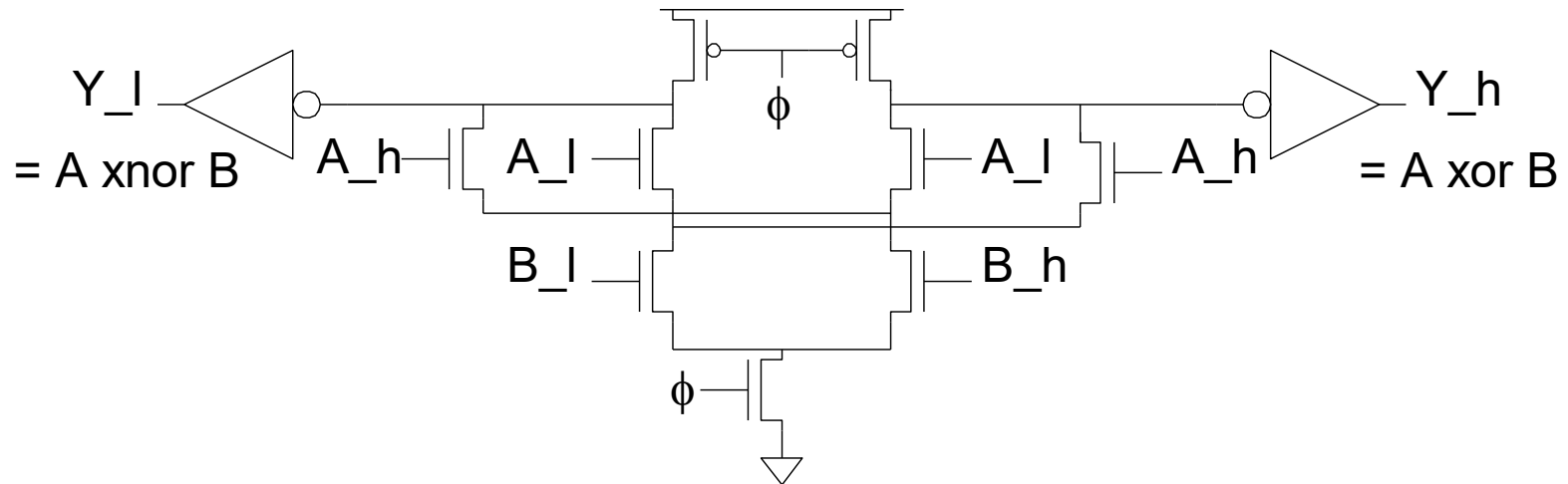
# Example: AND/NAND

- ❑ Given  $A_h, A_l, B_h, B_l$
- ❑ Compute  $Y_h = AB, Y_l = \overline{AB}$
- ❑ Pulldown networks are conduction complements



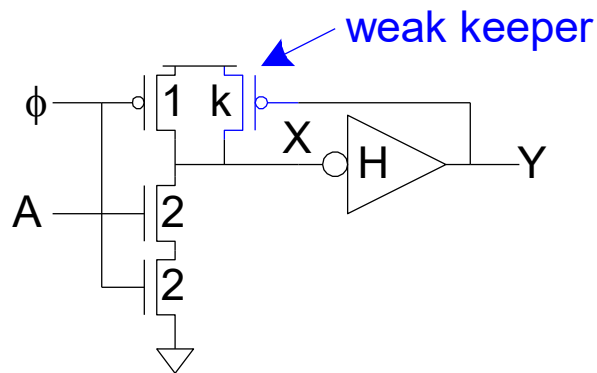
# Example: XOR/XNOR

- Sometimes possible to share transistors



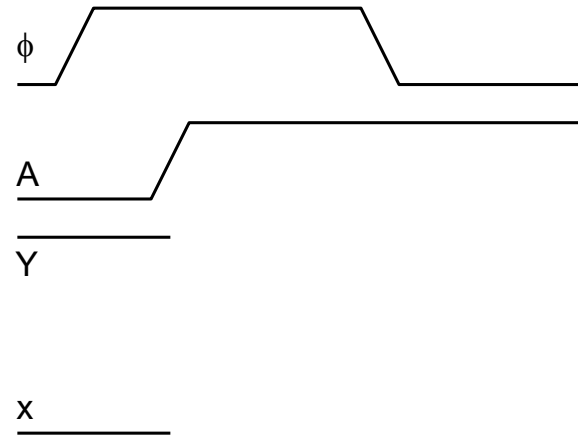
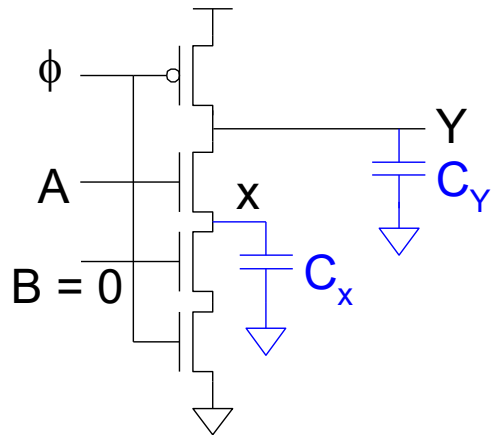
# Leakage

- ❑ Dynamic node floats high during evaluation
  - Transistors are leaky ( $I_{OFF} \neq 0$ )
  - Dynamic value will leak away over time
  - Formerly milliseconds, now nanoseconds
- ❑ Use keeper to hold dynamic node
  - Must be weak enough not to fight evaluation



# Charge Sharing

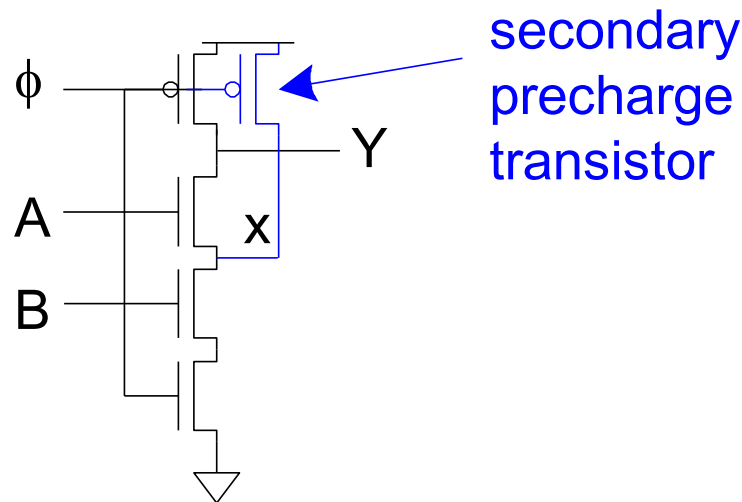
- ❑ Dynamic gates suffer from charge sharing



$$V_x = V_Y =$$

# Secondary Precharge

- ❑ Solution: add secondary precharge transistors
  - Typically need to precharge every other node
- ❑ Big load capacitance  $C_Y$  helps as well



# Noise Sensitivity

- ❑ Dynamic gates are very sensitive to noise
  - Inputs:  $V_{IH} \approx V_{tn}$
  - Outputs: floating output susceptible noise
- ❑ Noise sources
  - Capacitive crosstalk
  - Charge sharing
  - Power supply noise
  - Feedthrough noise
  - And more!

# Power

- ❑ Domino gates have high activity factors
  - Output evaluates and precharges
    - If output probability = 0.5,  $\alpha = 0.5$ 
      - Output rises and falls on half the cycles
  - Clocked transistors have  $\alpha = 1$
- ❑ Leads to very high power consumption

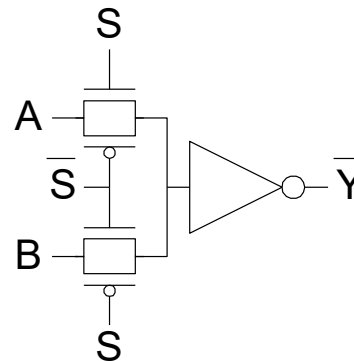
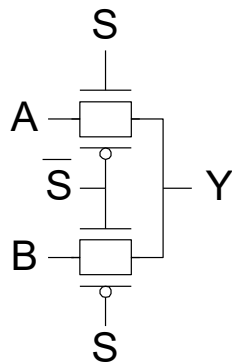


# Domino Summary

- ❑ Domino logic is attractive for high-speed circuits
  - 1.3 – 2x faster than static CMOS
  - But many challenges:
    - Monotonicity, leakage, charge sharing, noise
- ❑ Widely used in high-performance microprocessors in 1990s when speed was king
- ❑ Largely displaced by static CMOS now that power is the limiter
- ❑ Still used in memories for area efficiency

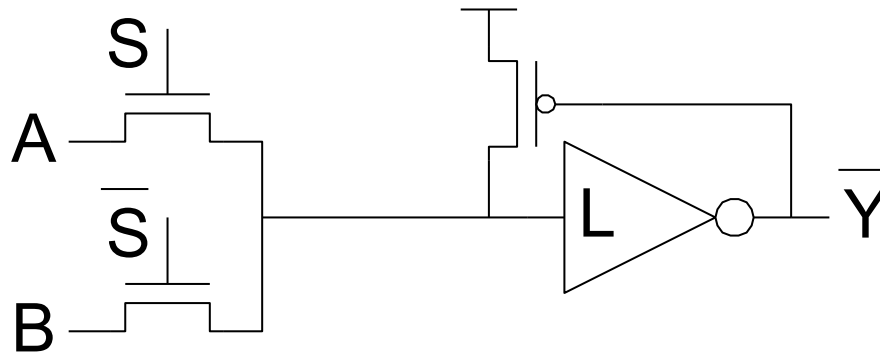
# Pass Transistor Circuits

- ❑ Use pass transistors like switches to do logic
- ❑ Inputs drive diffusion terminals as well as gates
- ❑ CMOS + Transmission Gates:
  - 2-input multiplexer
  - Gates should be restoring



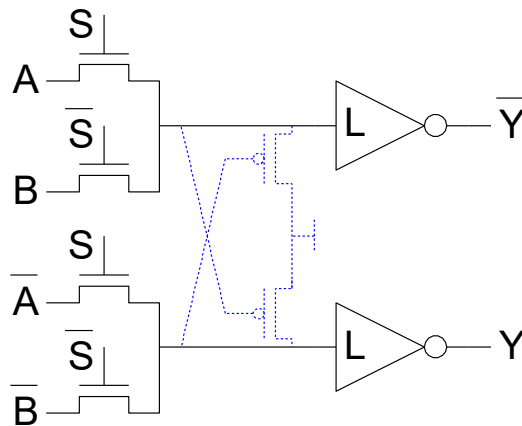
# LEAP

- ❑ **LEA**n integration with **P**ass transistors
- ❑ Get rid of pMOS transistors
  - Use weak pMOS feedback to pull fully high
  - Ratio constraint



# CPL

- ❑ Complementary Pass-transistor Logic
  - Dual-rail form of pass transistor logic
  - Avoids need for ratioed feedback
  - Optional cross-coupling for rail-to-rail swing



# Pass Transistor Summary

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- ❑ Researchers investigated pass transistor logic for general purpose applications in the 1990's
  - Benefits over static CMOS were small or negative
  - No longer generally used
- ❑ However, pass transistors still have a niche in special circuits such as memories where they offer small size and the threshold drops can be managed