



Department of Electrical Engineering
Indian Institute of Technology Kharagpur

Digital Signal Processing Laboratory (EE39203)
Autumn, 2022-23

Experiment 7: Fast Fourier Transform

Slot:

Date:

Student Name:

Roll No.:

Grading Rubric

	Tick the best applicable per row			Points
	Below Expectation	Lacking in Some	Meets all Expectation	
Completeness of the report				
Organization of the report (5 pts) <i>With cover sheet, answers are in the same order as questions in the lab, copies of the questions are included in report, prepared in LaTeX</i>				
Quality of figures (5 pts) <i>Correctly labelled with title, x-axis, y-axis, and name(s)</i>				
Understanding of the frequency range of DFT and effects of zero-padding (35 pts) <i>DFT and DTFT plots, Matlab code (DTFT samples), questions</i>				
Implementation of Divide-and-Conquer DFT and FFT (40 pts) <i>Matlab codes (dcDFT, fft2, fft4, fft8, fft stages), questions</i>				
Computation time comparison (15 pts) <i>Runtimes, questions</i>				
TOTAL (100 pts)				

Total Points (100):

TA Name:

TA Initials:

1. Introduction

This section continues the analysis of the DFT done earlier

$$DFT: X_N[k] = \sum_{n=0}^{N-1} x[n]e^{-j2\pi kn/N} \quad (17)$$

$$Inverse\ DFT: x[n] = \frac{1}{N} \sum_{k=0}^{N-1} X_N[k]e^{-j2\pi kn/N} \quad (18)$$

1.1. Shifting the Frequency Range

In this section, we will illustrate a representation for the DFT of equation (17) that is a bit more intuitive. First create a Hamming window x of length $N = 20$, using the Matlab command `x = hamming(20)`. Then use your matlab function `DFTsum` to compute the 20 point DFT of x . Plot the magnitude of the DFT, $|X_{20}[k]|$, versus the index k . Remember that the DFT index k starts at 0 not 1!

Lab Report:

Hand in the plot of the $|X_{20}[k]|$. Circle the regions of the plot corresponding to low frequency components.

Our plot of the DFT has two disadvantages. First, the DFT values are plotted against k rather than the frequency ω . Second, the arrangement of frequency samples in the DFT goes from 0 to 2π rather than from $-\pi$ to π , as is conventional with the DTFT. In order to plot the DFT values similar to a conventional DTFT plot, we must compute the vector of frequencies in radians per sample, and then “rotate” the plot to produce the more familiar range, $-\pi$ to π .

Let’s first consider the vector w of frequencies in radians per sample. Each element of w should be the frequency of the corresponding DFT sample $X[k]$, which can be computed by

$$\omega = 2\pi k/N \quad k \in [0, \dots, N - 1] \quad (19)$$

However, the frequencies should also lie in the range from $-\pi$ to π . Therefore, if $\omega \geq \pi$, then it should be set to $\omega - 2\pi$. An easy way of making this change in Matlab 5.1 is `w(w>=pi) = w(w>=pi)-2*pi`.

The resulting vectors X and w are correct, but out of order. To reorder them, we must swap the first and second halves of the vectors. Fortunately, Matlab provides a function specifically for this purpose, called `fftshift`.

Write a Matlab function to compute samples of the DTFT and their corresponding frequencies in the range $-\pi$ to π . Use the syntax

$$[X,w] = DTFTsamples(x)$$

where \mathbf{x} is an N point vector, \mathbf{X} is the length N vector of DTFT samples, and \mathbf{w} is the length N vector of corresponding radial frequencies. Your function *DTFTsamples* should call your function *DFTsum* and use the Matlab function *fftshift*.

Use your function *DTFTsamples* to compute DTFT samples of the Hamming window of length $N = 20$. Plot the magnitude of these DTFT samples versus frequency in rad/sample.

Lab Report:

1. Hand in the code for your function *DTFTsamples*.
2. Hand in the plot of the magnitude of the DTFT samples.

1.2. Zero Padding

The spacing between samples of the DTFT is determined by the number of points in the DFT. This can lead to surprising results when the number of samples is too small. In order to illustrate this effect, consider the finite-duration signal

$$x[n] = \begin{cases} \sin(0.1\pi n) & : 0 \leq n \leq 49 \\ 0 & : \text{otherwise} \end{cases} \quad (20)$$

In the following, you will compute the DTFT samples of $x[n]$ using both $N = 50$ and $N = 200$ point DFT's. Notice that when $N = 200$, most of the samples of $x[n]$ will be zeros because $x[n] = 0$ for $n \geq 50$. This technique is known as “zero padding”, and may be used to produce a finer sampling of the DTFT.

For $N = 50$ and $N = 200$, do the following:

1. Compute the vector \mathbf{x} containing the values $x[0], \dots, x[N - 1]$.
2. Compute the samples of $X[k]$ using your function *DTFTsamples*.
3. Plot the magnitude of the DTFT samples versus frequency in rad/sample.

Lab Report:

1. Submit your two plots of the DTFT samples for $N = 50$ and $N = 200$.
2. Which plot looks more like the true DTFT?
3. Explain why the plots look so different.

2. The Fast Fourier Transform Algorithm

We have seen in the preceding sections that the DFT is a very computationally intensive operation. In 1965, Cooley and Tukey¹ published an algorithm that could be used to compute the DFT much more efficiently. Various forms of their algorithm, which came to be known as the fast Fourier transform (FFT), had actually been developed much earlier by other mathematicians (even dating back to Gauss). It was their paper, however, which stimulated a revolution in the field of signal processing.

It is important to keep in mind at the outset that the FFT is not a new transform. It is simply a very efficient way to compute an existing transform, namely the DFT. As we saw, a straight forward implementation of the DFT can be computationally expensive

¹ J. W. Cooley and J. W. Tukey, “An algorithm for the machine calculation of complex Fourier series,” *Mathematics of Computation*, vol. 19, no. 90, p. 297-301, April 1965.

because the number of multiplies grows as the square of the input length (i.e. N^2 for an N point DFT). The FFT reduces this computation using two simple but important concepts. The first concept, known as divide-and-conquer, splits the problem into two smaller problems. The second concept, known as recursion, applies this divide-and-conquer method repeatedly until the problem is solved.

Consider the defining equation for the DFT and assume that N is even, so that $N/2$ is an integer:

$$X[k] = \sum_{n=0}^{N-1} x[n]e^{-j2\pi kn/N} \quad (21)$$

Here we have dropped the subscript of N in the notation for $X[k]$. We will also use the notation

$$X[k] = \text{DFT}_N\{x[n]\}$$

to denote the N point DFT of the signal $x[n]$. Suppose we break the sum in (21) into two sums, one containing all the terms for which n is even, and one containing all the terms for which n is odd:

$$X[k] = \sum_{\substack{n=0 \\ n \text{ even}}}^{N-1} x[n]e^{-j2\pi kn/N} + \sum_{\substack{n=0 \\ n \text{ odd}}}^{N-1} x[n]e^{-j2\pi kn/N} \quad (22)$$

We can eliminate the conditions “ n even” and “ n odd” in (22) by making a change of variable in each sum. In the first sum, we replace n by $2m$. Then as we sum m from 0 to $N/2 - 1$, $n = 2m$ will take on all even integer values between 0 and $N - 2$. Similarly, in the second sum, we replace n by $2m + 1$. Then as we sum m from 0 to $N/2 - 1$, $n = 2m + 1$ will take on all odd integer values between 0 and $N - 1$. Thus, we can write

$$X[k] = \sum_{m=0}^{N/2-1} x[2m]e^{-j2\pi k2m/N} + \sum_{m=0}^{N/2-1} x[2m+1]e^{-j2\pi k(2m+1)/N} \quad (23)$$

Next we rearrange the exponent of the complex exponential in the first sum, and split and rearrange the exponent in the second sum to yield

$$X[k] = \sum_{m=0}^{N/2-1} x[2m]e^{-j2\pi km/(N/2)} + e^{-j2\pi k/N} \sum_{m=0}^{N/2-1} x[2m+1]e^{-j2\pi km/(N/2)} \quad (24)$$

Now compare the first sum in (24) with the definition for the DFT given by (21). They have exactly the same form if we replace N everywhere in (21) by $N/2$. Thus the first sum in (8) is an $N/2$ point DFT of the even-numbered data points in the original sequence. Similarly, the second sum in (8) is an $N/2$ point DFT of the odd-numbered data points in the original sequence. To obtain the N point DFT of the complete sequence, we multiply the DFT of the odd-numbered data points by the complex exponential factor $e^{-j2\pi k/N}$, and then simply sum the two $N/2$ point DFTs.

To summarize, we will rewrite (24) according to this interpretation. First, we define two new $N/2$ point data sequences $x_0[n]$ and $x_1[n]$, which contain the even and odd-numbered data points from the original N point sequence, respectively:

$$x_0[n] = x[2n] \quad x_1[n] = x[2n + 1] \tag{25}$$

where $n = 0, \dots, N/2$. This separation of even and odd points is called *decimation* in time. The N point DFT of $x[n]$ is then given by

$$X[k] = X_0[k] + e^{-j2\pi k/N} X_1[k] \text{ for } k = 0, \dots, N - 1 \tag{26}$$

where $X_0[k]$ and $X_1[k]$ are the $N/2$ point DFT's of the even and odd points.

$$X_0[k] = \text{DFT}_{N/2}\{x_0[n]\} \quad X_1[k] = \text{DFT}_{N/2}\{x_1[n]\} \tag{27}$$

While equation (26) requires less computation than the original N point DFT, it can still be further simplified. First, note that each $N/2$ point DFT is periodic with period $N/2$. This means that we need to only compute $X_0[k]$ and $X_1[k]$ for $N/2$ values of k rather than the N values shown in (10). Furthermore, the complex exponential factor $e^{-j2\pi k/N}$ has the property that

$$-e^{-j2\pi \frac{k}{N}} = e^{-j2\pi \frac{k+N/2}{N}}$$

These two facts may be combined to yield a simpler expression for the N point DFT:

$$\left. \begin{aligned} X[k] &= X_0[k] + W_N^k X_1[k] \\ X[k + N/2] &= X_0[k] - W_N^k X_1[k] \end{aligned} \right\} \text{ for } k = 0, \dots, \frac{N}{2} - 1 \tag{28}$$

where the complex constants defined by $W_N^k = e^{-j2\pi k/N}$ are commonly known as the *twiddle factors*.

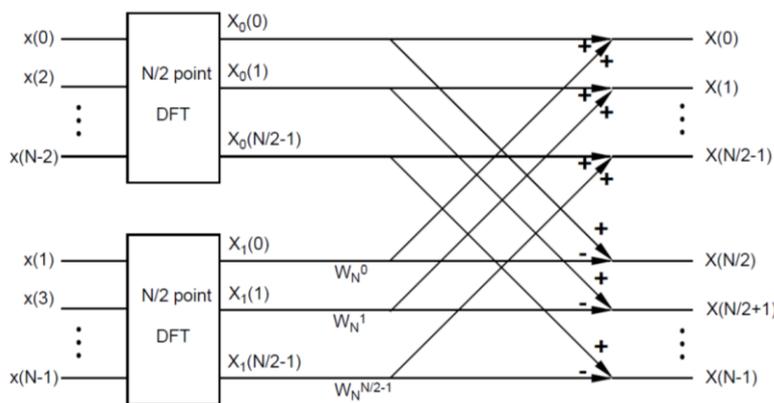


Figure 2: Divide and conquer DFT of equation (28). The N -point DFT is computed using the two $N/2$ -point DFT's $X_0^{(N/2)}[k]$ and $X_1^{(N/2)}[k]$.

Figure 2 shows a graphical interpretation of (28) which we will refer to as the “divide-and-conquer DFT”. We start on the left side with the data separated into even and odd subsets. We perform an $N/2$ point DFT on each subset, and then multiply the output of

the odd DFT by the required twiddle factors. The first half of the output is computed by adding the two branches, while the second half is formed by subtraction. This type of flow diagram is conventionally used to describe a fast Fourier transform algorithm.

2.1. Implementation of Divide-and-Conquer DFT

In this section, you will implement the DFT transformation using equation (28) and the illustration in Fig. 2. Write a Matlab function with the syntax

$$X = \text{dcDFT}(x)$$

where x is a vector of even length N , and X is its DFT. Your function `dcDFT` should do the following:

1. Separate the samples of x into even and odd points. Hint: The Matlab command `x0 = x(1:2:N)` can be used to obtain the “even” points.
2. Use your function `DFTsum` to compute the two $N/2$ point DFT’s.
3. Multiply by the twiddle factors $W_N^k = e^{-j2\pi k/N}$
4. Combine the two DFT’s to form X .

Test your function `dcDFT` by using it to compute the DFT’s of the following signals.

1. $x[n] = \delta[n]$ for $N = 10$
2. $x[n] = 1$ for $N = 10$
3. $x[n] = e^{j2\pi n/10}$ for $N = 10$

Lab Report:

Do the following:

1. Submit the code for your function `dcDFT`.
2. Determine the number of multiplies that are required in this approach to computing an N point DFT. (Consider a multiply to be one multiplication of real or complex numbers.) HINT: Refer to the diagram of Fig. 2, and remember to consider the $N/2$ point DFTs.

2.2. Recursive Divide and Conquer

The second basic concept underlying the FFT algorithm is that of recursion. Suppose $N/2$ is also even. Then we may apply the same decimation-in-time idea to the computation of each of the $N/2$ point DFT’s in Fig. 2. This yields the process depicted in Fig. 3. We now have two stages of twiddle factors instead of one.

How many times can we repeat the process of decimating the input sequence? Suppose N is a power of 2, i.e. $N = 2^p$ for some integer p . We can then repeatedly decimate the sequence until each subsequence contains only two points. It is easily seen from (21) that the 2 point DFT is a simple sum and difference of values.

$$\begin{aligned} X[0] &= x[0] + x[1] \\ X[1] &= x[0] - x[1] \end{aligned} \tag{29}$$

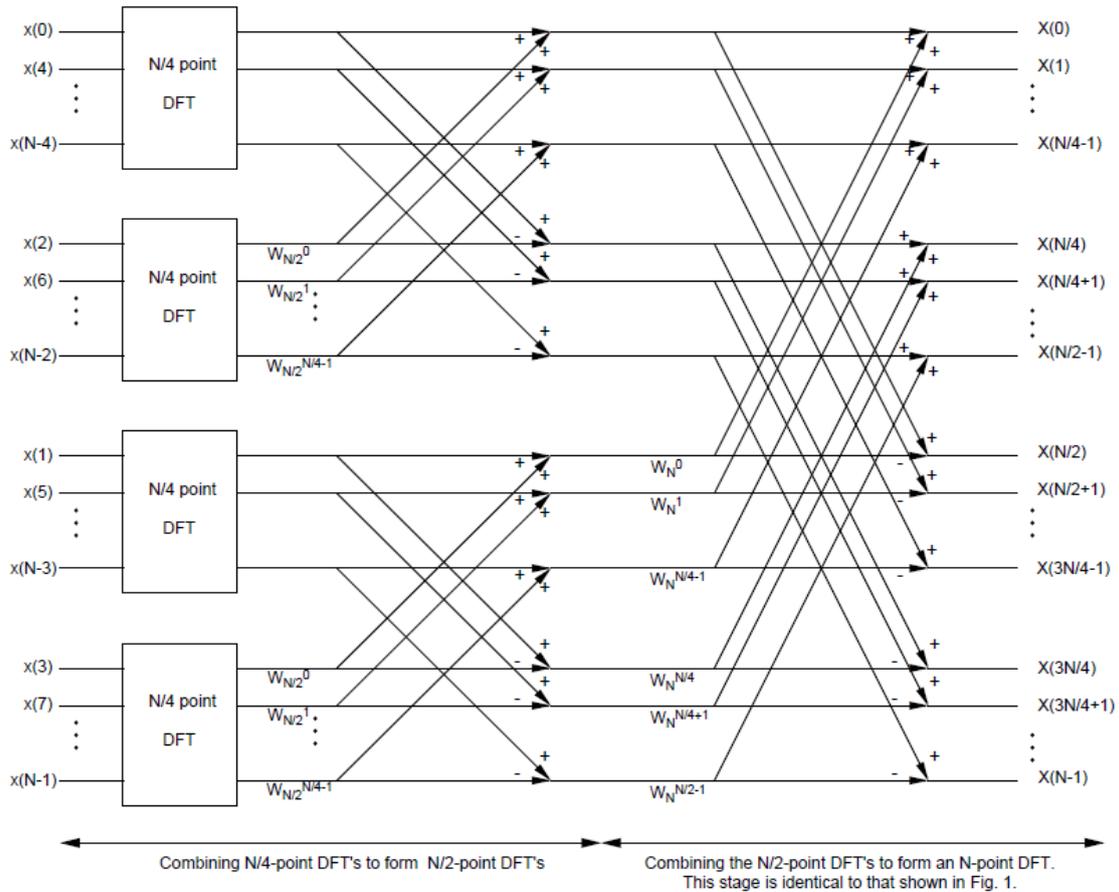


Figure 3: Recursion of the decimation-in-time process. Now each $N/2$ -point is calculated by combining two $N/4$ -point DFT's.

Fig. 4 shows the flow diagram that results for an 8 point DFT when we decimate 3 times. Note that there are 3 stages of twiddle factors (in the first stage, the twiddle factors simplify to "1"). This is the flow diagram for the complete decimation-in-time 8 point FFT algorithm. How many multiplies are required to compute it?

Write three Matlab functions to compute the 2, 4, and 8-point FFT's using the syntax

$$X = \text{FFT2}(x)$$

$$X = \text{FFT4}(x)$$

$$X = \text{FFT8}(x)$$

The function FFT2 should directly compute the 2-point DFT using (29), but the functions FFT4 and FFT8 should compute their respective FFT's using the divide and conquer strategy. This means that FFT8 should call FFT4 , and FFT4 should call FFT2 .

Test your function FFT8 by using it to compute the DFT's of the following signals. Compare these results to the previous ones.

1. $x[n] = \delta[n]$ for $N = 8$
2. $x[n] = 1$ for $N = 8$
3. $x[n] = e^{j2\pi n/8}$ for $N = 8$

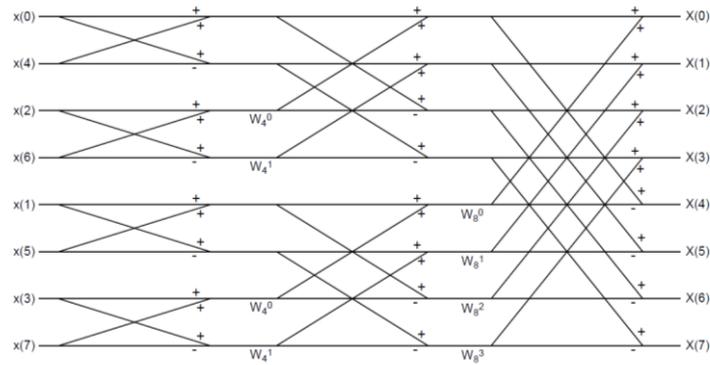


Figure 4: 8-point FFT

Lab Report:

1. Submit the code for your functions *FFT2*, *FFT4* and *FFT8*.
2. List the output of *FFT8* for the case $x[n] = 1$ for $N = 8$.
3. Calculate the total number of multiplies by twiddle factors required for your 8-point FFT. (A multiply is a multiplication by a real or complex number.)
4. Determine a formula for the number of multiplies required for an $N = 2p$ point FFT. Leave the expression in terms of N and p . How does this compare to the number of multiplies required for direct implementation when $p = 10$?

If you wrote the *FFT4* and *FFT8* functions properly, they should have almost the exact same form. The only difference between them is the length of the input signal, and the function called to compute the $(N/2)$ -pt DFTs. Obviously, it's redundant to write a separate function for each specific length DFT when they each have the same form. The preferred method is to write a recursive function, which means that the function calls itself within the body. It is imperative that a recursive function has a condition for exiting without calling itself, otherwise it would never terminate.

Write a recursive function $X = \text{fft_stage}(x)$ that performs one stage of the FFT algorithm for a power-of-2 length signal. An outline of the function is as follows:

1. Determine the length of the input signal.
2. If $N = 2$, then the function should just compute the 2-pt DFT as in equation (29), and then return.
3. If $N > 2$, then the function should perform the FFT steps described previously (i.e. decimate, compute $(N/2)$ -pt DFTs, re-combine), calling *fft_stage* to compute the $(N/2)$ -pt DFTs.

Note that the body of this function should look very similar to previous functions written in this lab. Test *fft_stage* on the three 8-point signals given above, and verify that it returns the same results as *FFT8*.

Lab Report:

Submit the code for your *fft_stage* function.

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