

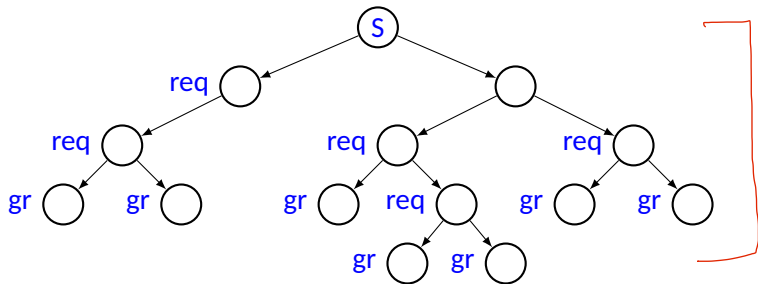
Artificial Intelligence: Foundations & Applications

Temporal Logic and Applications



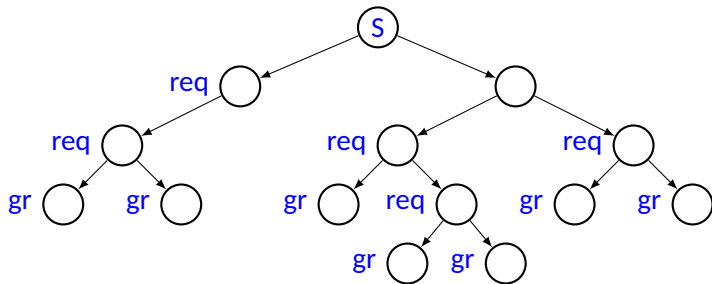
Prof. Partha P. Chakrabarti & Arijit Mondal
Indian Institute of Technology Kharagpur

CTL example



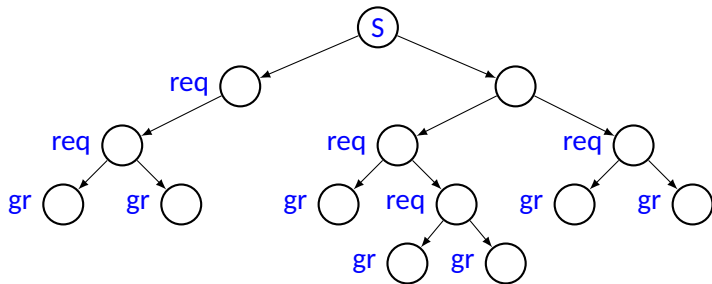
- ✓ • From S the system always makes a request in future: $A F req$

CTL example



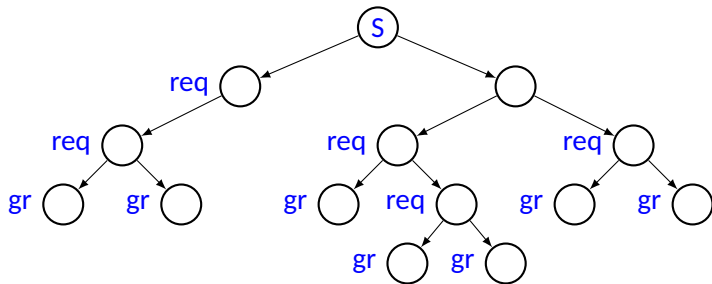
- From S the system always makes a request in future: $AFreq$

CTL example



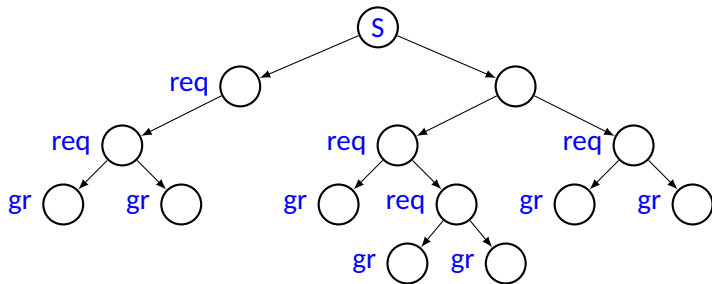
- From S the system always makes a request in future: $AF req$
- All requests are eventually granted: $AG(req \rightarrow AF gr)$

CTL example



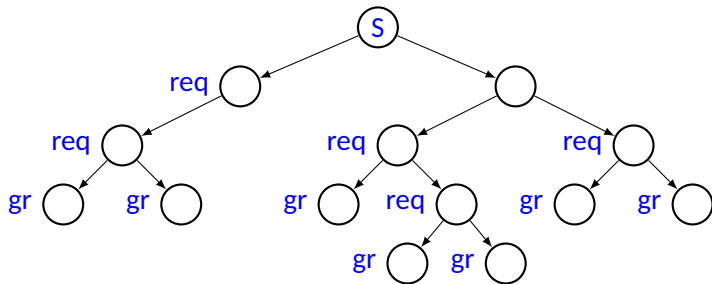
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CTL example



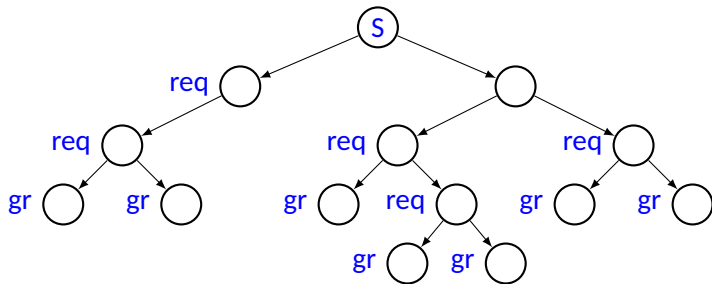
- From S the system always makes a request in future: $AF req$
- All requests are eventually granted: $AG(req \rightarrow AF gr)$
- Sometimes requests are immediately granted: $EF(\underline{req} \rightarrow \underline{EX gr})$

CTL example



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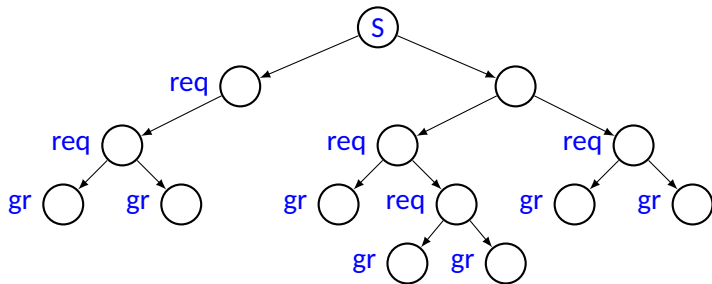
CTL example



- From S the system always makes a request in future: $AF req$
- All requests are eventually granted: $AG(req \rightarrow AF gr)$
- Sometimes requests are immediately granted: $EF(req \rightarrow EX gr)$
- Requests are held till grant is received:

$$AG(req \rightarrow A (req \vee gr))$$

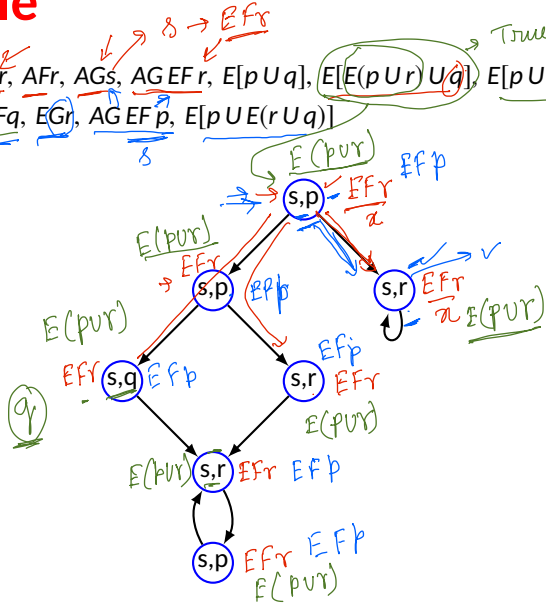
CTL example



- From S the system always makes a request in future: $AF req$
- All requests are eventually granted: $AG(req \rightarrow AF gr)$
- Sometimes requests are immediately granted: $EF(req \rightarrow EX gr)$
- Requests are held till grant is received: $AG(req \rightarrow A(req U gr))$

CTL example

- True properties: $\underline{EFr}$, $\underline{AFr}$, $\underline{AGs}$, $\underline{AGEFr}$, $\underline{E[p \cup q]}$, $\underline{E[E(p \cup r) \cup q]}$, $\underline{E[p \cup E(q \cup r)]}$
- False properties: $\underline{AFq}$, $\underline{EGr}$, $\underline{AGEFp}$, $\underline{E[p \cup E(r \cup q)]}$



CTL Model Checking

- It checks whether a given CTL formula f holds on a given Kripke structure M i.e., $M \models f$

- Need to have modalities for EX, EU and EG

- Other modalities can be expressed using EX, EU and EG

- $AF f \equiv \neg EG \neg f$
- $AG f \equiv \neg EF \neg f$
- $A(f U g) \equiv (\neg EG \neg g) \wedge (\neg E[\neg g U (\neg f \wedge \neg g)])$

- Basic procedure

- The set $Sat(f)$ of all states satisfying f is computed recursively
- $M \models f$ if and only if $S_0 \subseteq Sat(f)$

models

$M \models f$

AU	AG	AF	AX
EU	EG	EF	EX

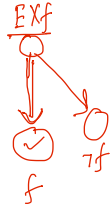
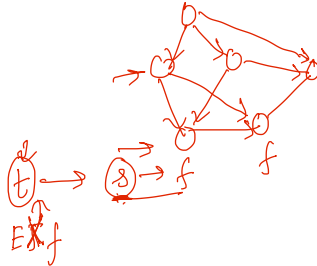
$f \models \text{true} \vee f$

CTL Model Checking: EX f

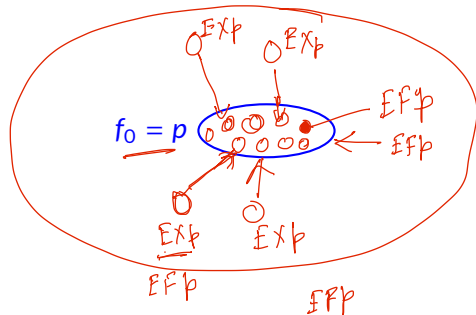
function CheckEX(f)

1. $S_f = \{s \in S \mid f \in L(s)\}$
2. while $S_f \neq \emptyset$
3. Choose $s \in S_f$
4. $S_f = S_f - \{s\}$
5. for all t such that $(t, s) \in T$
6. if $f \notin L(t)$
7. $L(t) = L(t) \cup \{EXf\}$
8. endif
9. end for
10. end while

$$T \equiv (PS, NS)$$



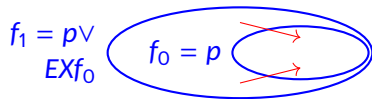
CTL Model Checking: $EFfp$



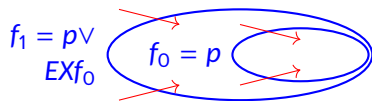
CTL Model Checking: $EF f$



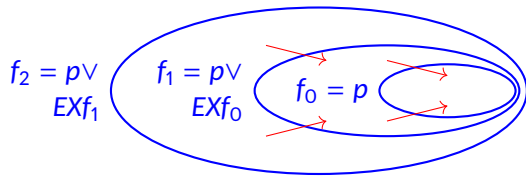
CTL Model Checking: $EF f$



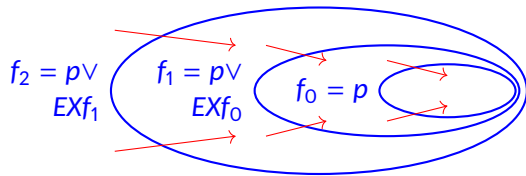
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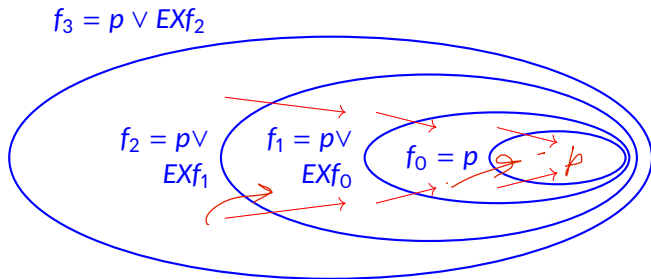
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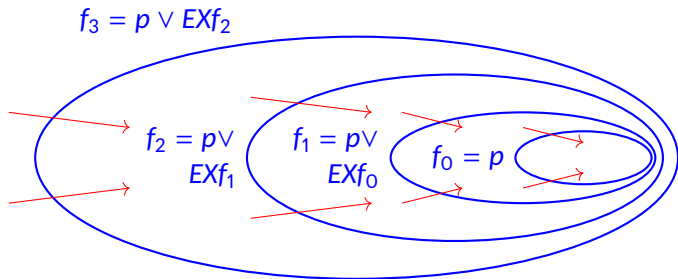
CTL Model Checking: $EF f$



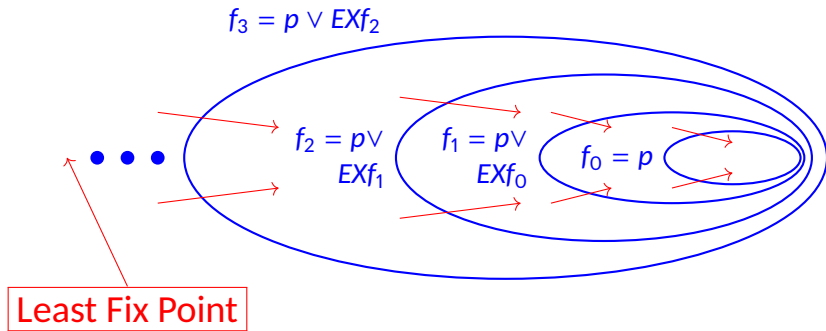
CTL Model Checking: $EF f$



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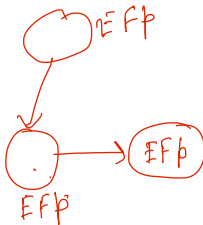
CTL Model Checking: $EF f$



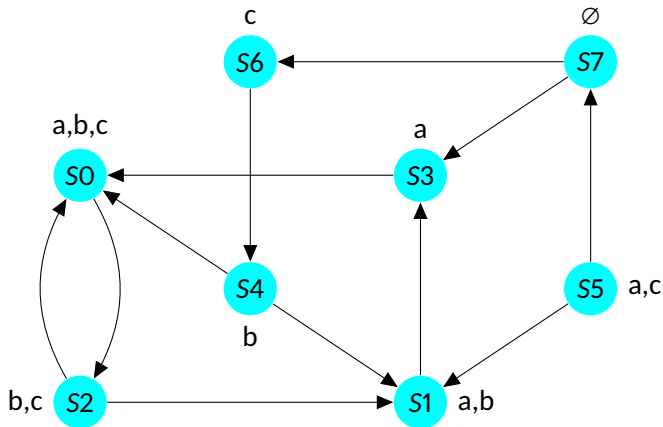
CTL Model Checking: $EF\ p$

function CheckEF(p)

1. $S_p = \{s \in S \mid p \in L(s)\}$
2. for all $s \in S_p$ do $L(s) = L(s) \cup \{EFp\}$
3. while $S_p \neq \emptyset$
4. Choose $s \in S_p$
5. $S_p = S_p - \{s\}$
6. for all t such that $(t, s) \in T$
7. if $\{EFp\} \notin L(t)$
8. $L(t) = L(t) \cup \{EFp\}$
9. $S_p = S_p \cup \{t\}$
10. endif
11. end for
12. end while



Example: $g = \underline{EF((a \oplus c) \wedge (a \oplus b))}$

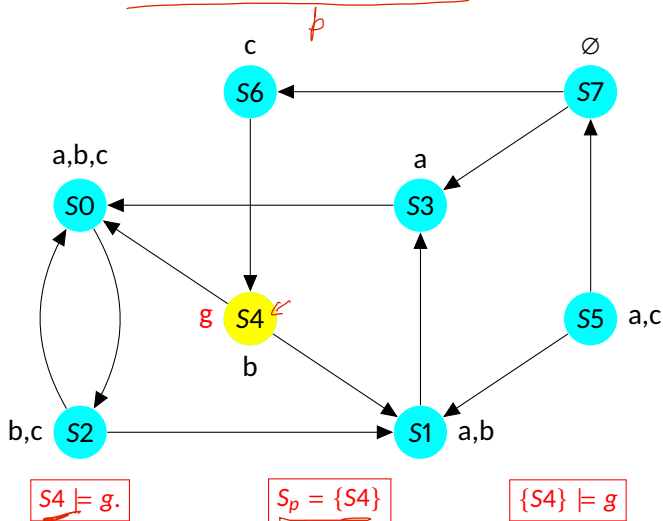


Let $p = \underline{((a \oplus c) \wedge (a \oplus b))}$

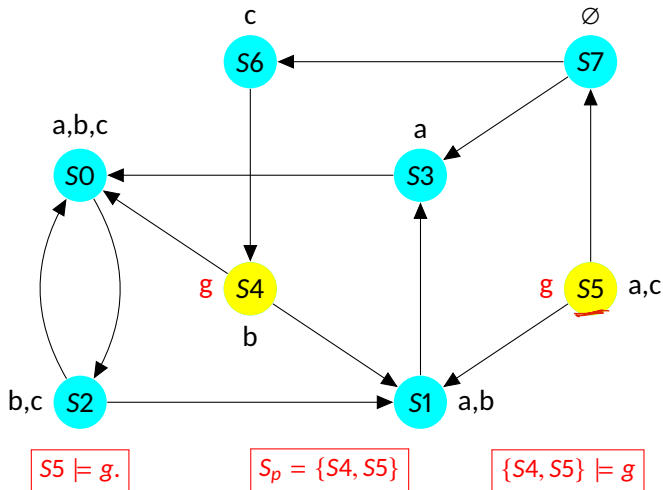
$S_p = \{\}$

$\{\} \models \underline{g}$

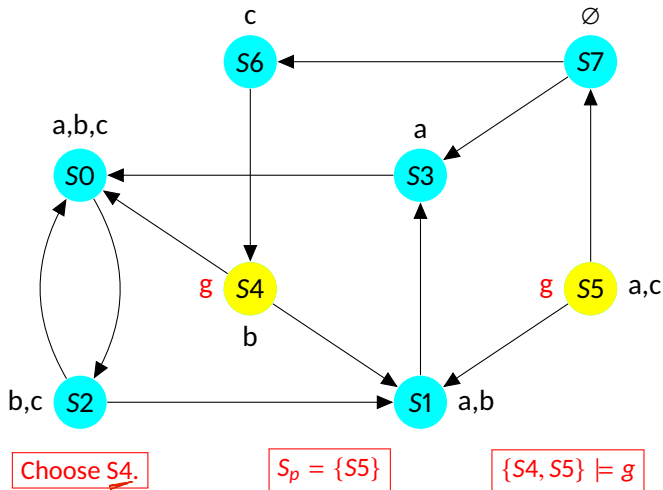
Example: $g = EF((a \oplus c) \wedge (a \oplus b))$



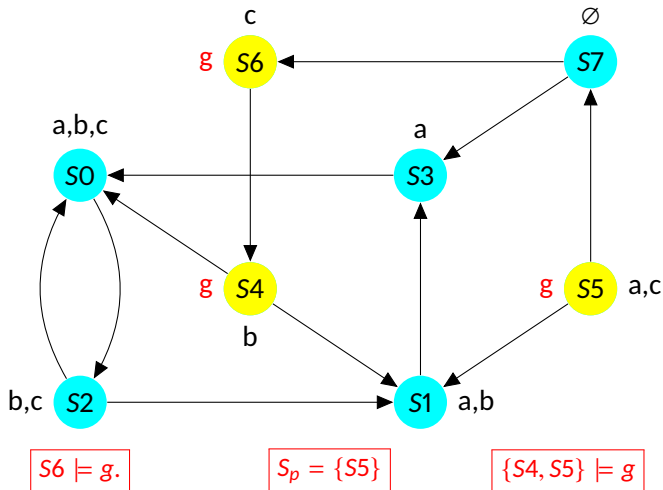
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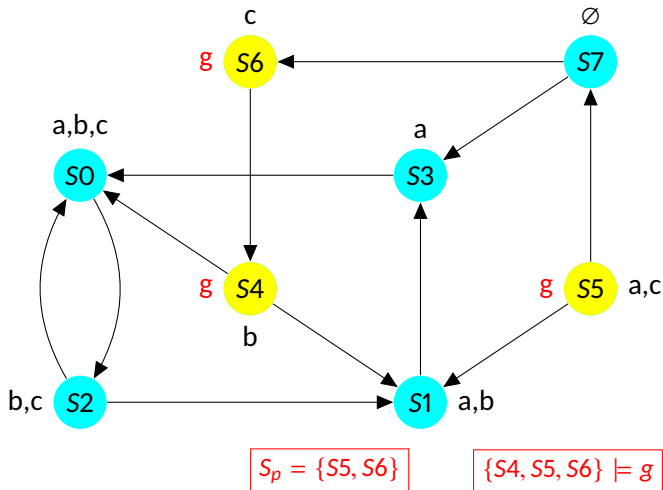
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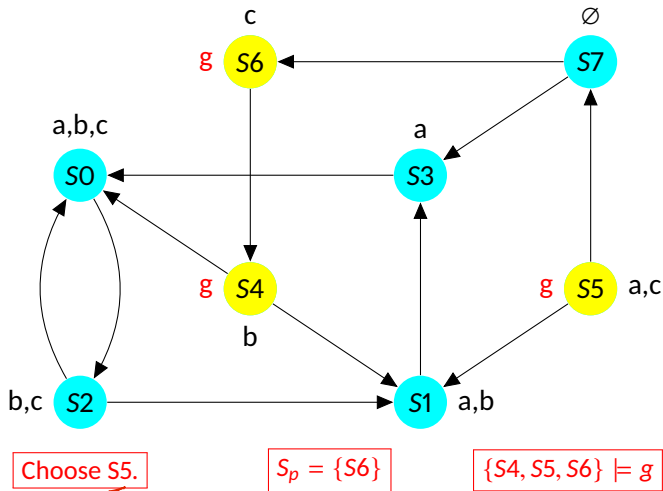
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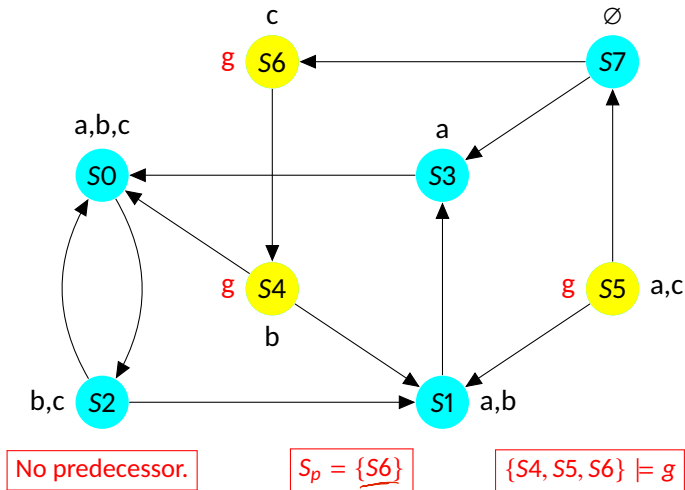
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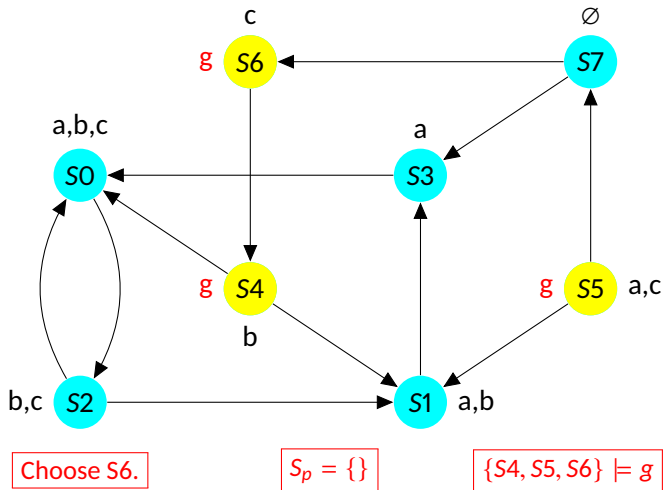
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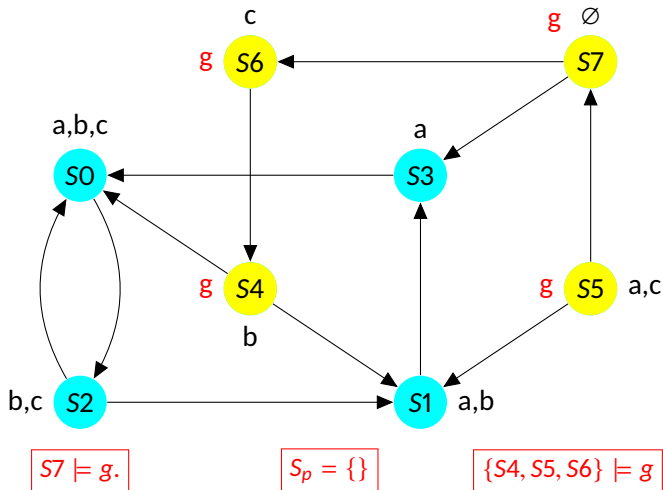
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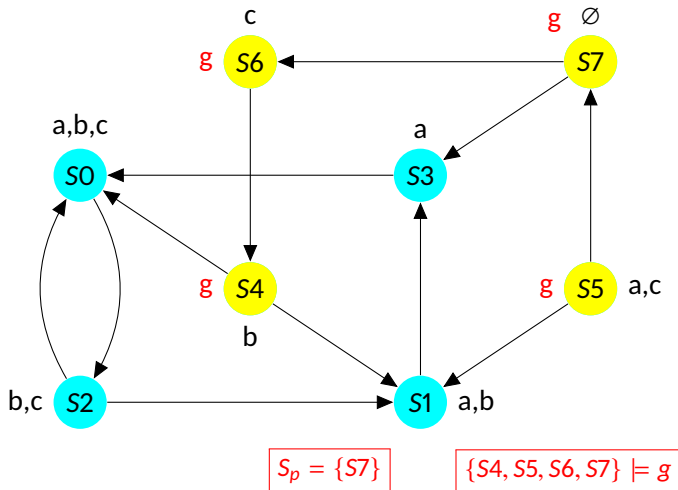
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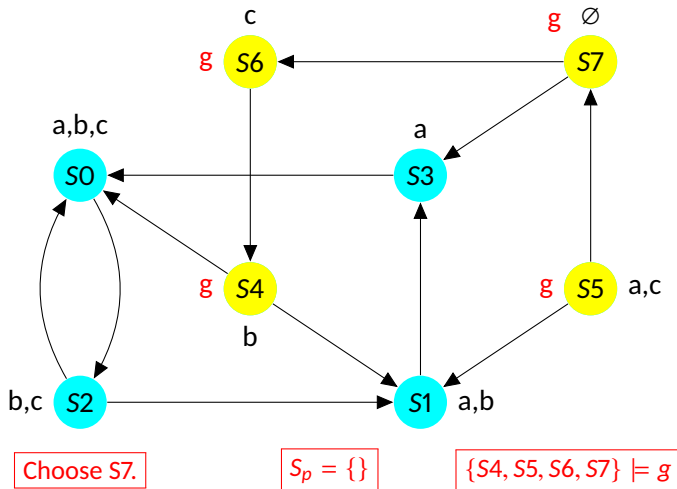
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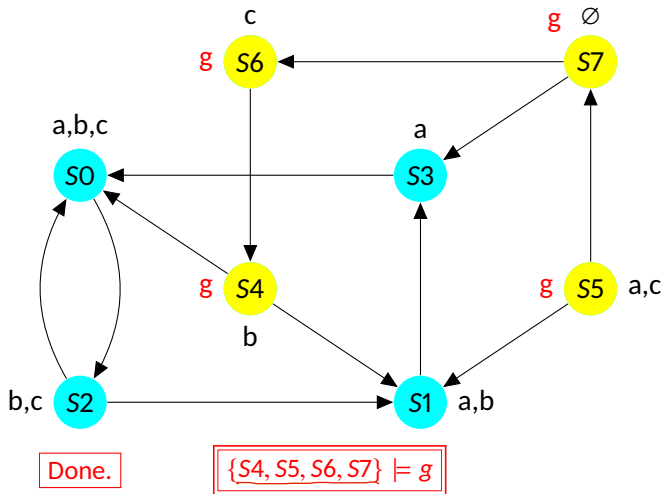
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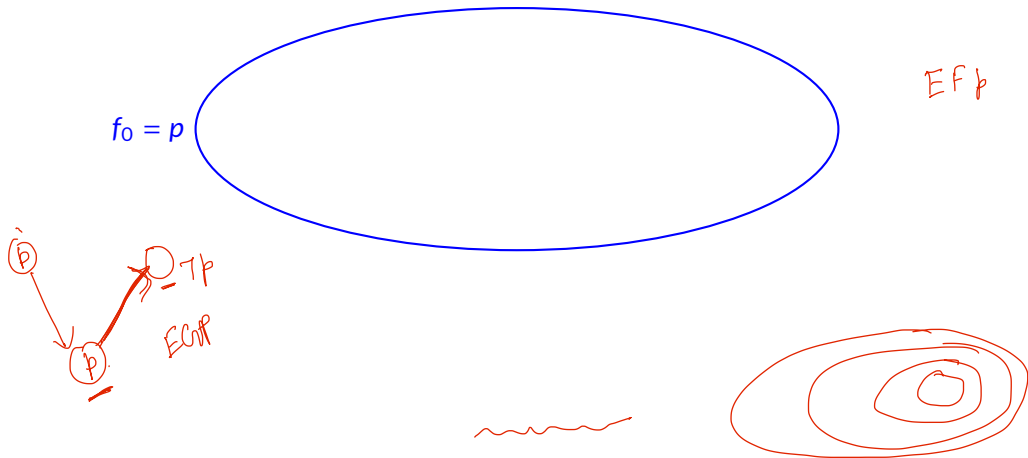
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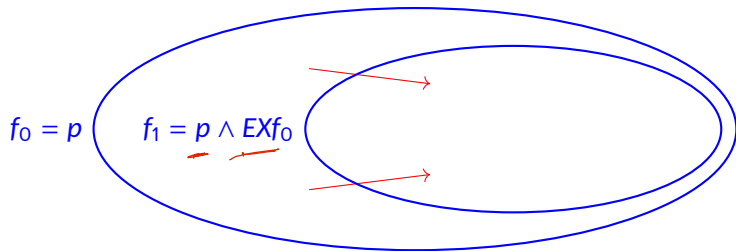
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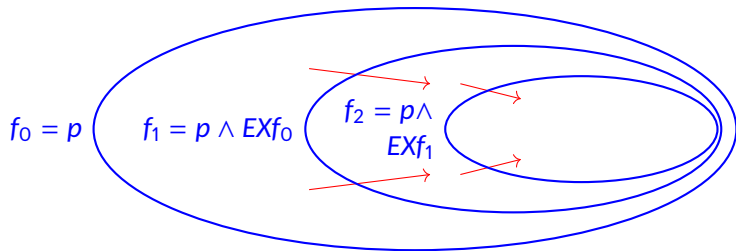
CTL Model Checking: $EG\ p$



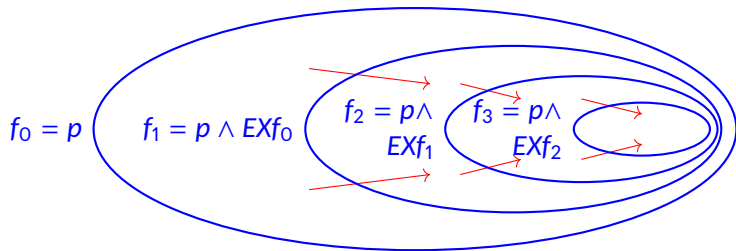
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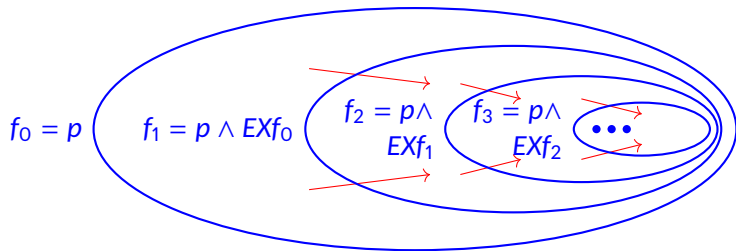
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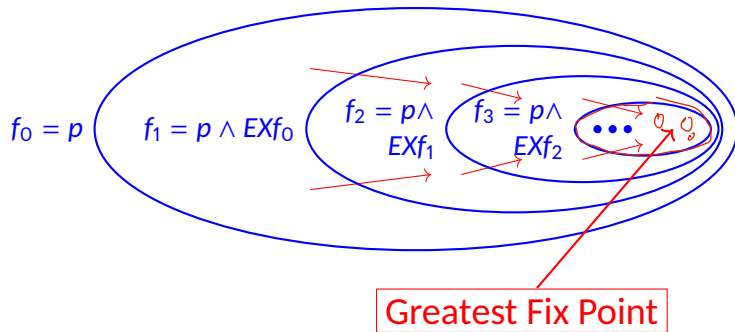
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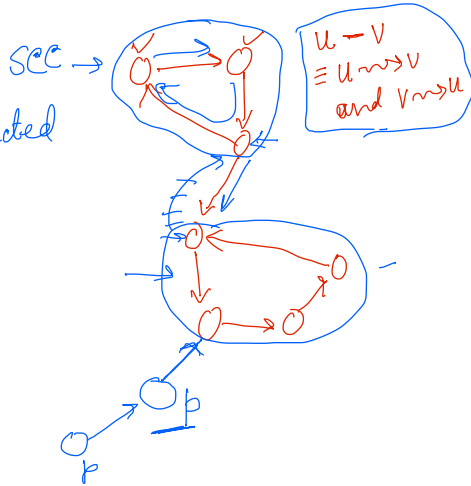


CTL Model Checking: EGp

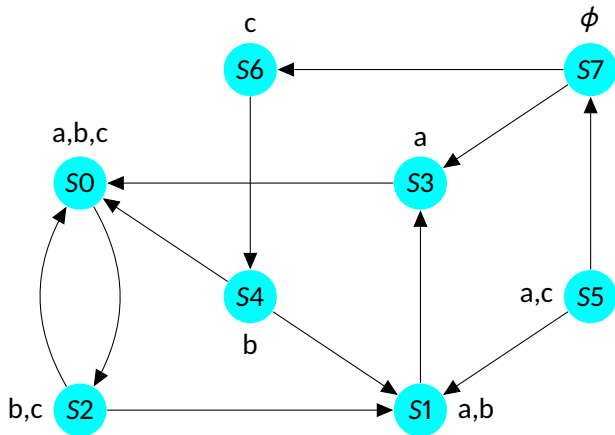
function CheckEG(p)

1. $S_p = \{s \in S \mid p \in L(s)\}$
2. $\text{SCC} = \{C \mid C \text{ is nontrivial SCC of } S_p\}$
3. $R = \bigcup_{C \in \text{SCC}} \{s \mid s \in C\}$
4. for all $s \in R$ do $L(s) = L(s) \cup \{EGp\}$
5. while $R \neq \emptyset$
6. Choose $s \in R$
7. $R = R - \{s\}$
8. for all t such that $(t, s) \in T$ and $t \in S_p$
9. if $\{EGp\} \notin L(t)$
10. $L(t) = L(t) \cup \{EGp\}$
11. $R = R \cup \{t\}$
12. endif
13. end for
14. end while

Strongly Connected
Component

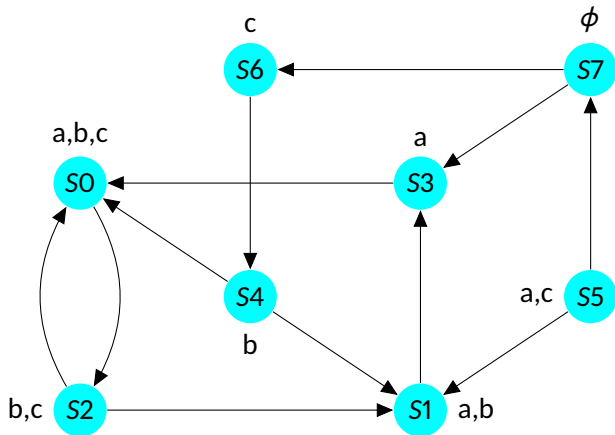


Example: $g = EG\ b$



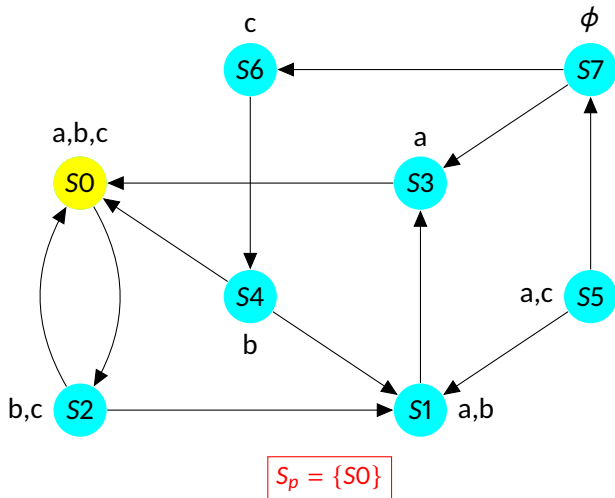
Let $p = b$

Example: $g = EG\ b$

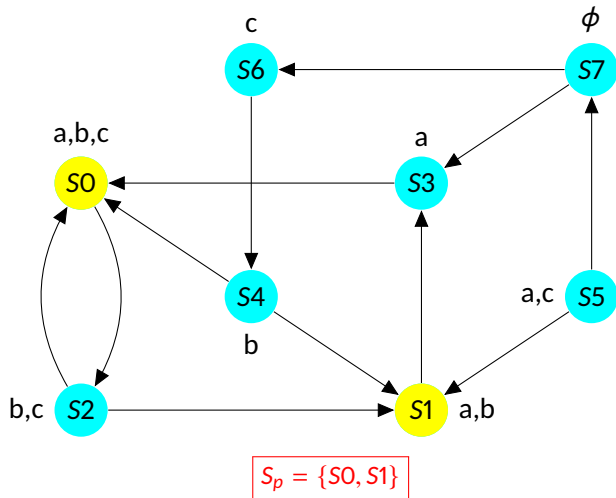


Find states satisfying b .

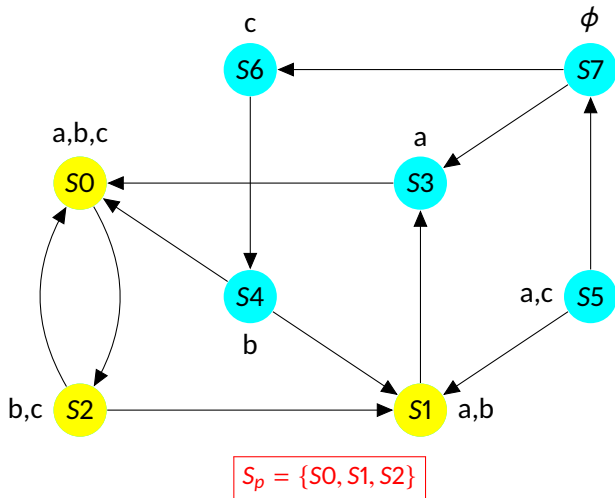
Example: $g = EG\ b$



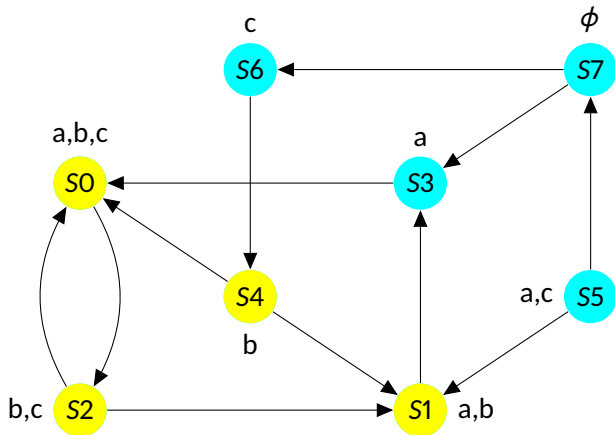
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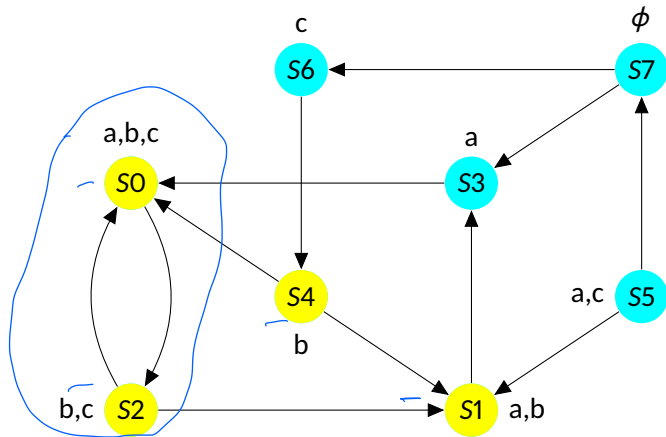


Example: $g = EG\ b$



$$S_p = \{S_0, S_1, S_2, S_4\}$$

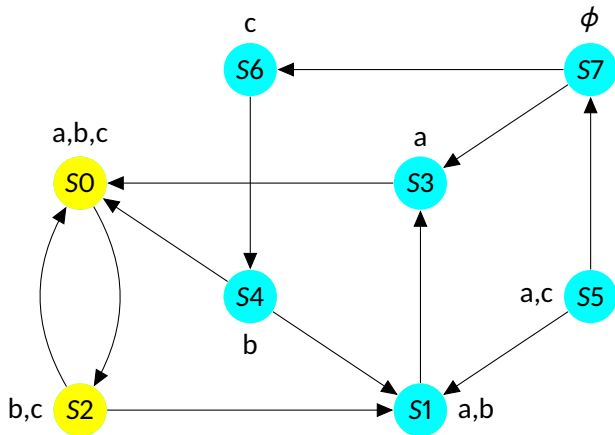
Example: $g = EG\ b$



Find SCC using S_p .

$S_p = \{S_0, S_1, S_2, S_4\}$

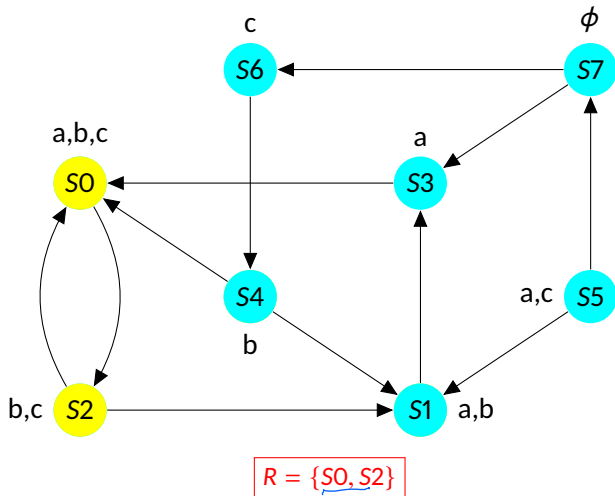
Example: $g = EG\ b$



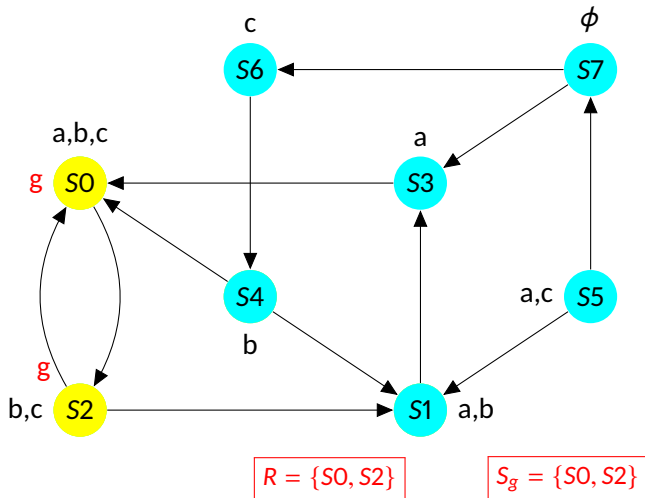
Find SCC using S_p .

$S_p = \{S0, S1, S2, S4\}$

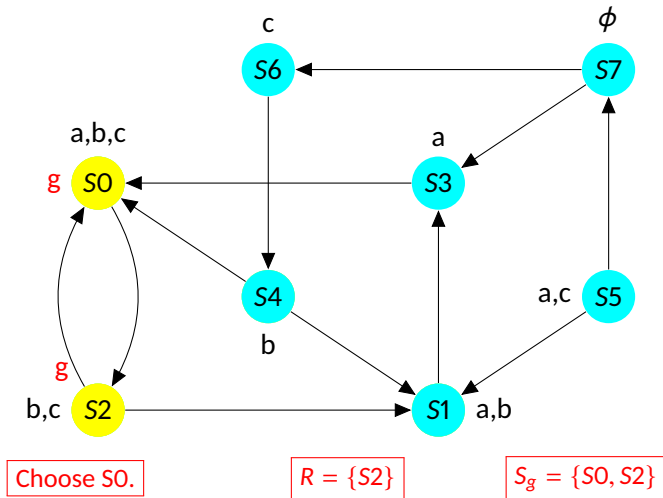
Example: $g = EG\ b$



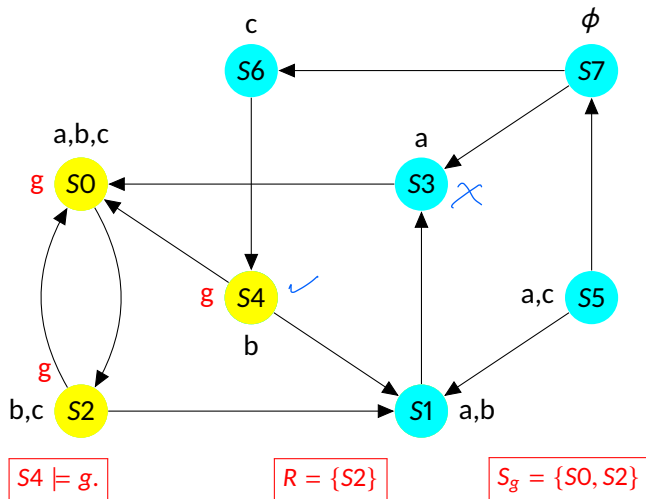
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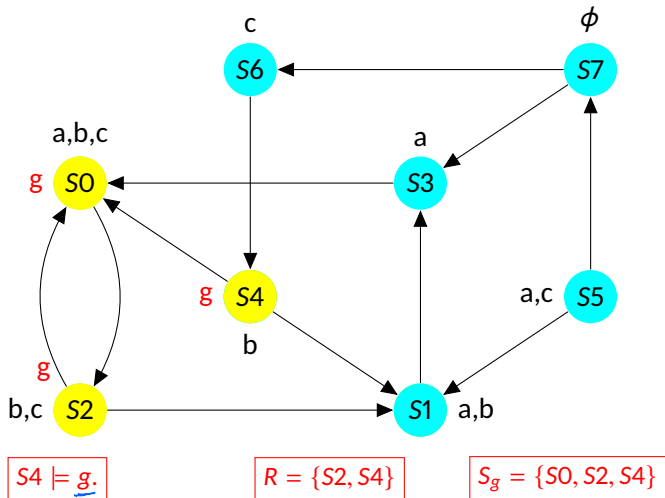
Example: $g = EG\ b$



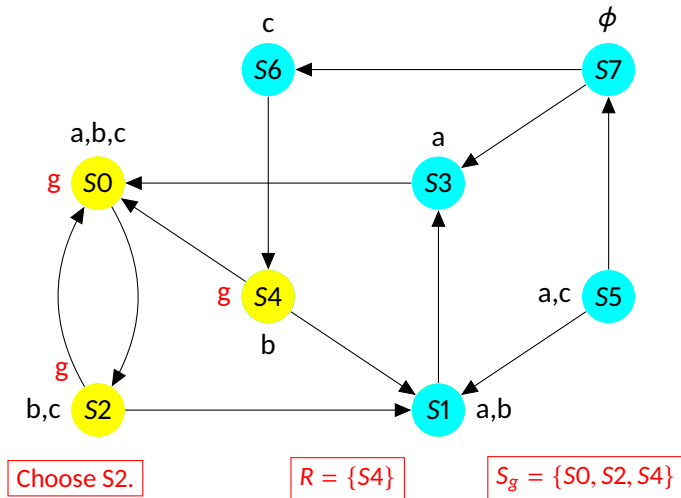
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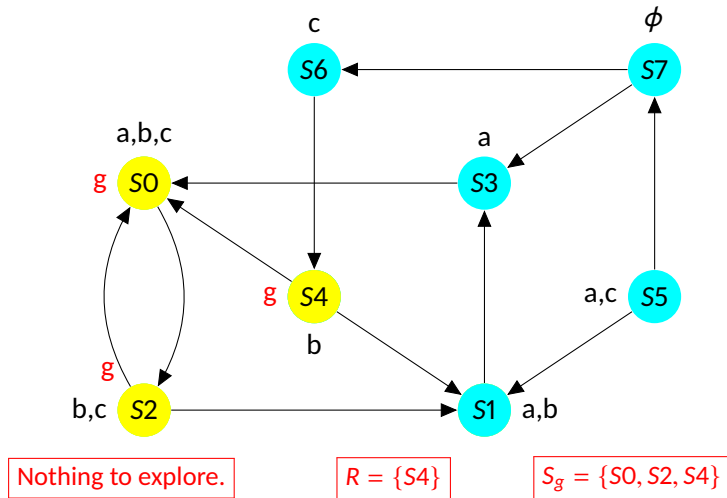
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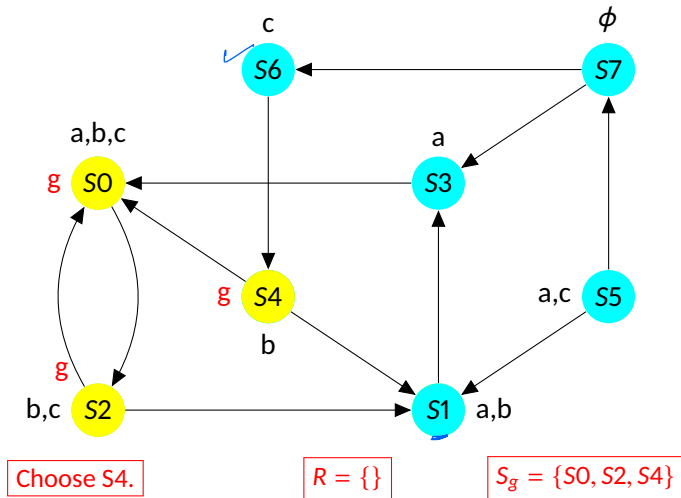
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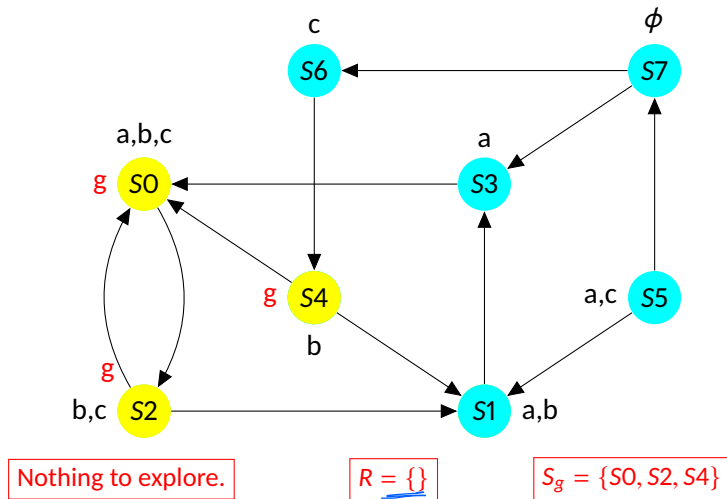
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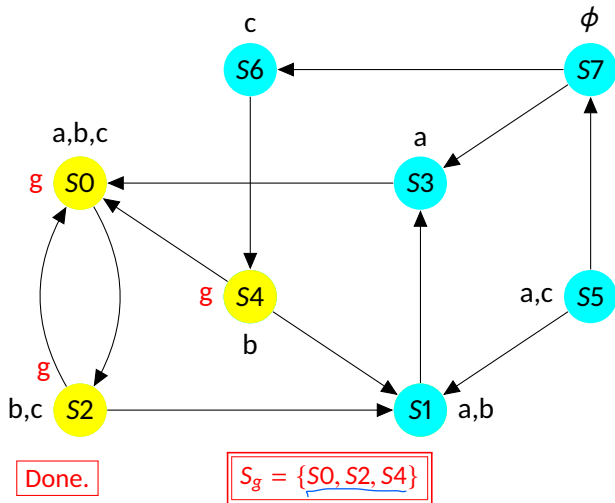
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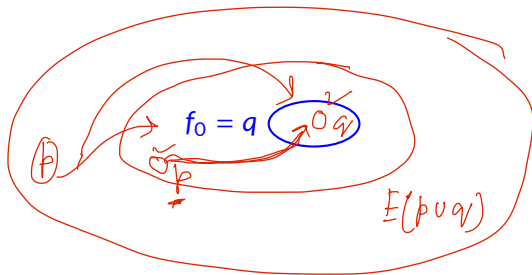
Example: $g = EG\ b$



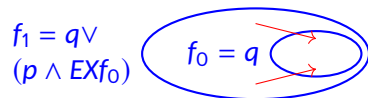
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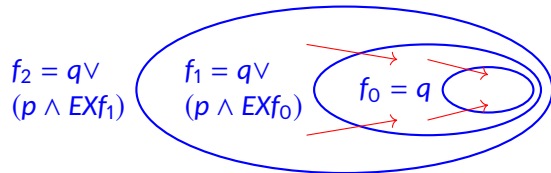
CTL Model Checking: $E(p \text{ U } q)$



CTL Model Checking: $E(p \ U \ q)$

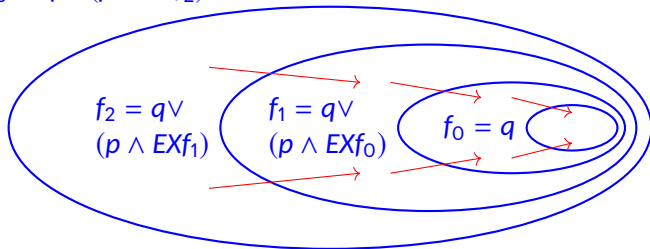


CTL Model Checking: $E(p \ U \ q)$

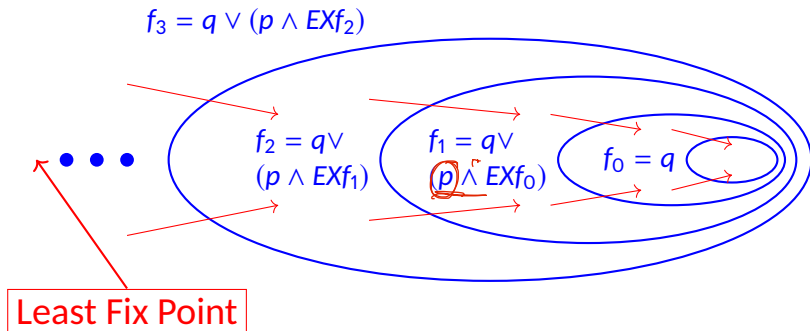


CTL Model Checking: $E(p \cup q)$

$$f_3 = q \vee (p \wedge EXf_2)$$



CTL Model Checking: $E(p \ U \ q)$

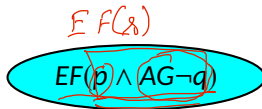


CTL Model Checking: $E(\underline{p} U q)$

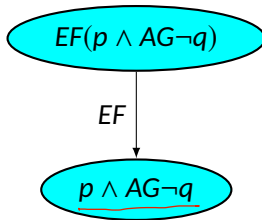
function CheckEU(p,q)

1. $S_q = \{s \in S \mid q \in L(s)\}$
2. **for all** $s \in S_q$ **do** $L(s) = L(s) \cup \{E(p U q)\}$
3. **while** $S_q \neq \emptyset$
4. **Choose** $s \in S_q$
5. $S_q = S_q - \{s\}$
6. **for all** t **such that** $(t, s) \in T$
7. **if** $\{E(p U q)\} \notin L(t)$ **and** $p \in L(t)$
8. $L(t) = L(t) \cup \{E(p U q)\}$
9. $S_q = S_q \cup \{t\}$
10. **endif**
11. **end for**
12. **end while**

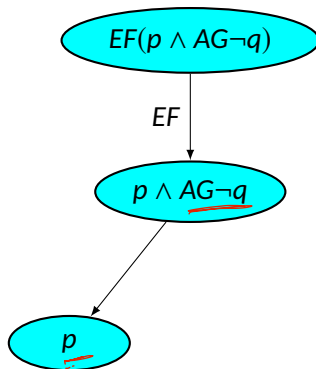
Nested CTL query

$$EF(\varphi)$$

$$EF(p \wedge AG\neg q)$$

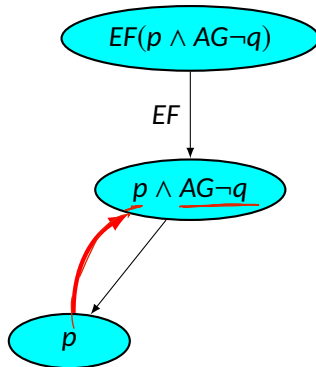
Nested CTL query



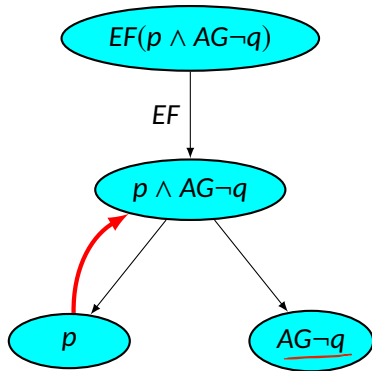
Nested CTL query



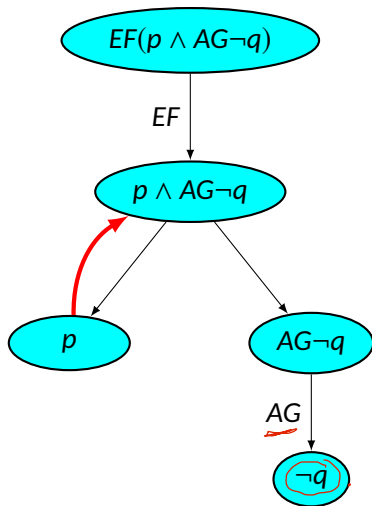
Nested CTL query



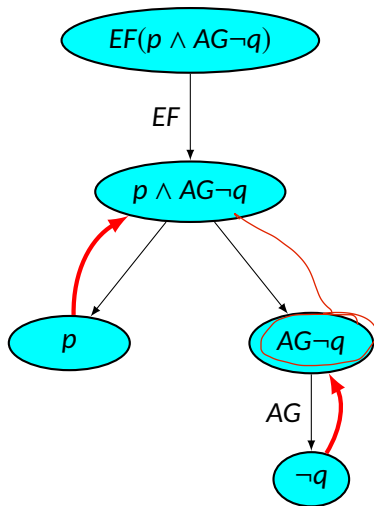
Nested CTL query



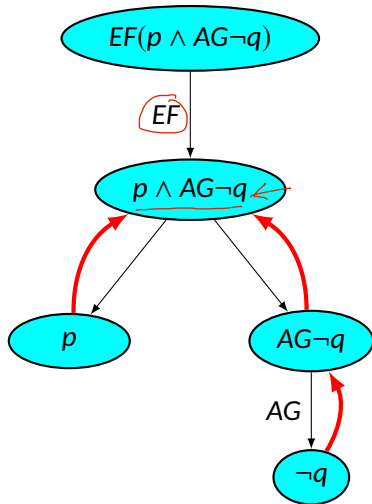
Nested CTL query



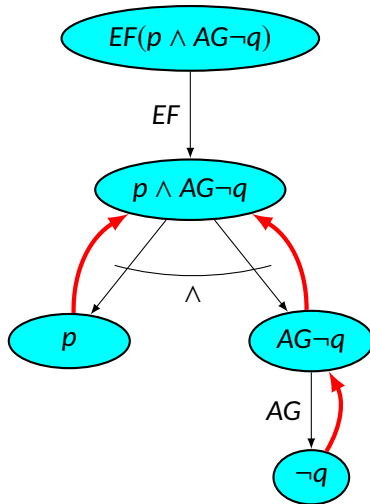
Nested CTL query



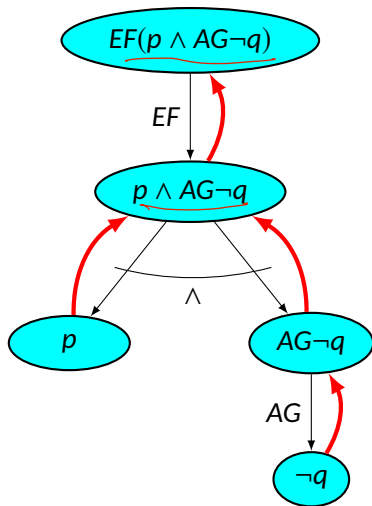
Nested CTL query



Nested CTL query



Nested CTL query



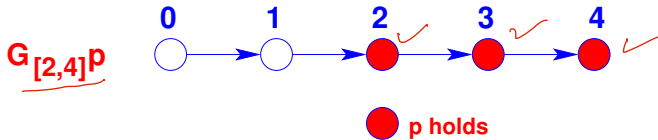
Real Time properties

- Real time systems
 - Predictable response time are necessary for correct operation
 - Safety critical systems like controller for aircraft, industrial machinery are a few examples
- It is difficult to express complex timing properties
 - Simple: “event p will happen in future”
 - Fp
 - Complex: “event p will happen within at most n time units”
 - $p \vee (Xp) \vee (XXp) \vee \dots ([XX \dots n \text{ times}]p)$

Bounded Temporal Operators

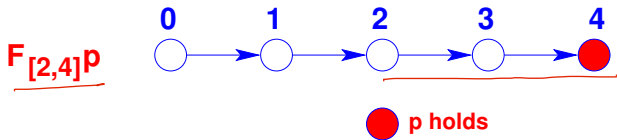
- Specify real-time constraints
 - Over bounded traces
- Various bounded temporal operators
 - $G_{[m,n]}p$ - p always holds between m^{th} and n^{th} time step
 - $F_{[m,n]}p$ - p eventually holds between m^{th} and n^{th} time step
 - $X_{[m]}p$ - p holds at the m^{th} time step
 - $p U_{[m,n]} q$ - q eventually holds between m^{th} and n^{th} time step and p holds until that point of time

Examples



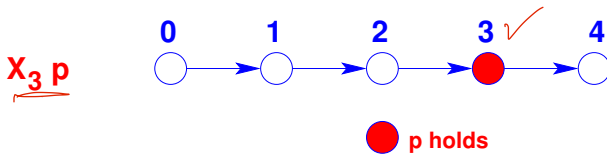
- **p** holds always between 2nd and 4th time step

Examples



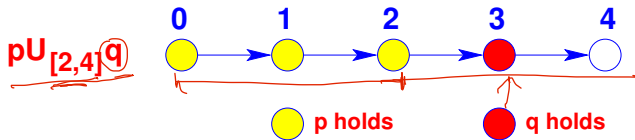
- **p** holds eventually between 2nd and 4th time step

Examples



- p holds in the 3rd time step

Examples



- **q** holds eventually between 2nd and 4th time step and **p** holds until **q** holds

Timing properties

- Whenever request is recorded grant should take place within 4 time units
 - $AG(\text{posedge}(\text{req}) \rightarrow \text{AF}_{[0,4]} \text{posedge}(\text{gr}))$
- The arbiter will provide exactly 64 time units to high priority user in each grant
 - $AG(\text{posedge}(\text{hpusing}) \rightarrow \text{A}(\neg \text{negedge}(\text{hpusing}) \text{ U}_{[64,64]} \text{negedge}(\text{hpusing})))$

Thank you!