

Artificial Intelligence: Foundations & Applications

SAT solvers



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Introduction

- SAT is one of the central problems in computer science community that has both theoretical as well as practical challenges
- This was the first NP-Complete problem }
- Wide variety of application domains - formal verification, test pattern generation, planning, scheduling, time tabling, etc.
- Provides a generic framework for combinatorial reasoning and search platform |
- It is based upon propositional logic (Boolean logic)
- CSP problems can be mapped to SAT
- There exists good open-source industrial strength SAT solvers


SAT problems


- Propositions - $\mathcal{P} = \{\underline{a}, \underline{b}, \underline{c}, \dots\}$!
- Literals - $\{\underline{a}, \underline{\neg a}, b, \neg b, \dots\}$
- Clause - $C_1 = (\underline{a} \vee \underline{b} \vee \neg c)$, $C_2 = (\neg a \vee b \vee \neg d), \dots$
 - Clause is disjunction of literals !
- Formula - $\mathcal{F} = C_1 \wedge C_2 \wedge \dots$
 - Conjunctive normal form (CNF) !
- Goal is to find an assignment (interpretation) to the propositions such that \mathcal{F} is true
 - \mathcal{F} is satisfiable if there exists at least one valid interpretation
 - \mathcal{F} is unsatisfiable if there exists none

$$f = \underbrace{(a \wedge b \wedge c)}_{\substack{\downarrow \\ 1 \quad 1 \quad 1}} \vee \neg a \wedge \neg b \wedge c$$

SAT tools

- Very good open-source SAT solvers are available

- 
- MiniSAT
 - zChaff
 - CaDiCaL
 - Glucose
 - Lingeling

- 
- PicoSAT
 - Cryptominisat
 - Rsat
 - Riss
 - many others

- <http://www.satcompetition.org/> ←

Input format - DIMACS

- There is standard format to specify clauses and its literals
- To specify comments

c This line is comment

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- To specify problem, you need to provide number of variables and number of clauses

→ c p cnf num_of_variables num_of_clauses
p cnf 3 4

Input format - DIMACS

- There is standard format to specify clauses and its literals

- To specify comments

`c This line is comment`

- To specify problem, you need to provide number of variables and number of clauses

`c p cnf num_of_variables num_of_clauses`

`p cnf 4 4`

- To specify CNF

`c list_of_literals 0`

→ `1 -2 3 0`

→ `2 4 0` → $b \vee d$

→ `-3 0` → $\neg c$

`-1 2 3 -4 0`

$\neg a \vee b \vee c \vee \neg d$

a, b, c
 $\neg a, \neg b, \neg c$
 $1 \quad 2 \quad 3$
 $-1 \quad -2 \quad -3$

$a \vee \neg b \vee c$
 $1 \quad -2 \quad 3$

Output format

- Outputs from a SAT solver are - SATISFIABLE / UNSATISFIABLE, an assignment of Boolean variables
- Typically it will be as follows

→ SAT

→ -1 2 -3 4 0 ←

- The last line needs to be interpreted as follows: $\neg a \wedge b \wedge \neg c \wedge d$
- There may be additional messages to provide information on resource usage, statistics, etc.

SAT modeling: Propositional logic - 1

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 - If \mathcal{M} is tautology then $\mathcal{F} \wedge \bar{\mathcal{G}}$ will be **false** ie. $\mathcal{F} \wedge \mathcal{G} = \emptyset$ ✓
 - If \mathcal{M} is satisfiable then so is $\mathcal{F} \wedge \bar{\mathcal{G}}$

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$$\left\{ \begin{array}{ll} \text{p} & \text{cnf} \quad 2 \quad 3 \\ -1 & 2 \quad 0 \\ 1 & 0 \\ -2 & 0 \end{array} \right.$$

UNSATISFIABLE

$$\boxed{(\neg a \vee b) \wedge a} \wedge \neg b$$
$$\underline{b \wedge a \wedge \neg b} = \phi$$

SAT modeling: Propositional logic - 2

- If Rajat is the Director then Rajat is well known. Rajat is not the Director. So, Rajat is not well known.
- Propositions: a : Rajat is the Director, b : Rajat is well known.
- Formula (\mathcal{F}): $a \implies b$, $\neg a$
- Goal (\mathcal{G}): $\neg b$

SAT modeling: Propositional logic - 2

- If Rajat is the Director then Rajat is well known. Rajat is not the Director. So, Rajat is not well known.
- Propositions: a : Rajat is the Director, b : Rajat is well known.
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- Goal (\mathcal{G}): $\neg b$
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- Formula (\mathcal{F}): $a \implies b, \neg a$
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-1 2 0
-1 0
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- Propositions: a : Rajat is the Director, b : Rajat is well known.
- Formula (\mathcal{F}): $a \implies b, \neg a$
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```
p cnf 2 3
-1 2 0
-1 0
2 0
```

SATISFIABLE

SAT modeling: Propositional logic - 3

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- Target: Can we prove that the unicorn is mythical? Magical? Horned?
- Propositions: a -mythical, b -mortal, c -mammal, d -horned, e -magical

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 - $\neg a \implies (b \wedge c)$

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 - $(\neg b \vee c) \implies d$

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 - $d \implies e$

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$$\begin{array}{l} p \text{ cnf } 5 \ 7 \\ \left| \begin{array}{lll} -1 & -2 & 0 \\ 1 & 2 & 0 \\ 1 & 3 & 0 \end{array} \right. \end{array}$$

$$\left| \begin{array}{lll} 2 & 4 & 0 \\ -3 & 4 & 0 \\ -4 & 5 & 0 \end{array} \right.$$

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p cnf 5 7
-1 -2 0
1 2 0
1 3 0

2 4 0
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p cnf 5 7
-1 -2 0
1 2 0
1 3 0

2 4 0
-3 4 0
-4 5 0
-1 0 //

a - SAT

-5 0 // e - UNSAT

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p cnf 5 7
-1 -2 0
1 2 0
1 3 0

2 4 0
-3 4 0
-4 5 0
-1 0 // a - SAT

-5 0 // e - UNSAT
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p cnf 5 7
-1 -2 0
1 2 0
1 3 0

2 4 0
-3 4 0
-4 5 0
-1 0 //

a - SAT

-5 0 // e - UNSAT
-4 0 // d - UNSAT

SAT modeling: Propositional logic - 4

- Knowledge base: Holmes owns two suits: one black and one tweed. He always wears either a tweed suit or sandals. Whenever he wears his tweed suit and a purple shirt, he chooses to not wear a tie. He never wears the tweed suit unless he is also wearing either a purple shirt or sandals. Whenever he wears sandals, he also wears a purple shirt. Yesterday, Holmes wore a bow tie.
- Target: What else did he wear?
- Propositions: 1-*bs*-black suit, 2-*ts*-tweed suit, 3-*s*-sandals, 4-*ps*-purple shirt, 5-*t*-tie

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- Target: What else did he wear?
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- Formula (\mathcal{F}):

SAT modeling: Propositional logic - 4

- Knowledge base: Holmes owns two suits: one black and one tweed. He always wears either a tweed suit or sandals. Whenever he wears his tweed suit and a purple shirt, he chooses to not wear a tie. He never wears the tweed suit unless he is also wearing either a purple shirt or sandals. Whenever he wears sandals, he also wears a purple shirt. Yesterday, Holmes wore a bow tie.
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- Formula (\mathcal{F}):
 - $\neg bs \vee \neg ts$

SAT modeling: Propositional logic - 4

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 - $ts \vee s$

SAT modeling: Propositional logic - 4

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 - $ts \vee s$
 - $(ts \wedge ps) \implies \neg t$

SAT modeling: Propositional logic - 4

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SAT modeling: Propositional logic - 4

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SAT modeling: Propositional logic - 4

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 - $ts \implies (ps \vee s) \equiv (\neg ts \vee ps \vee s)$ • $s \implies ps$

SAT modeling: Propositional logic - 4

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SAT modeling: Propositional logic - 4

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SAT modeling: Propositional logic - 4

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- Goal (\mathcal{G}): All satisfying solutions

SAT modeling: Propositional logic - 4

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- SAT modeling

SAT modeling: Propositional logic - 4

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- Target: What else did he wear?

- Propositions: 1-*bs*-black suit, 2-*ts*-tweed suit, 3-*s*-sandals, 4-*ps*-purple shirt, 5-*t*-tie

- Formula (\mathcal{F}):

$$\begin{aligned} & \bullet \neg bs \vee \neg ts & \bullet ts \vee s & \bullet (ts \wedge ps) \implies \neg t \equiv (\neg ts \vee \neg ps \vee \neg t) \\ & \bullet ts \implies (ps \vee s) \equiv (\neg ts \vee ps \vee s) & \bullet s \implies ps \equiv (\neg s \vee ps) & \bullet t \end{aligned}$$

- Goal (\mathcal{G}): All satisfying solutions

- SAT modeling

$$\begin{array}{ll} p \text{ cnf } 5 \ 6 & -2 \ 4 \ 3 \ 0 \\ -1 \ -2 \ 0 & -3 \ 4 \ 0 \\ 2 \ 3 \ 0 & 5 \ 0 \\ -2 \ -4 \ -5 \ 0 & \end{array}$$

SAT modeling: Propositional logic - 4

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- $ts \implies (ps \vee s) \equiv (\neg ts \vee ps \vee s)$ • $s \implies ps \equiv (\neg s \vee ps)$ • t

- Goal (\mathcal{G}): All satisfying solutions

- SAT modeling

p cnf 5 6
-1 -2 0
2 3 0
-2 -4 -5 0

-2 4 3 0

-3 4 0

5 0

SAT: -1 -2 3 4 5 0

SAT modeling: Propositional logic - 4

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- SAT modeling

p cnf 5 6
-1 -2 0
2 3 0
-2 -4 -5 0

-2 4 3 0
-3 4 0
5 0

SAT: -1 -2 3 4 5 0

Add: 1 2 -3 -4 -5 0

SAT modeling: Propositional logic - 4

- Knowledge base: Holmes owns two suits: one black and one tweed. He always wears either a tweed suit or sandals. Whenever he wears his tweed suit and a purple shirt, he chooses to not wear a tie. He never wears the tweed suit unless he is also wearing either a purple shirt or sandals. Whenever he wears sandals, he also wears a purple shirt. Yesterday, Holmes wore a bow tie.

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- $ts \implies (ps \vee s) \equiv (\neg ts \vee ps \vee s)$ • $s \implies ps \equiv (\neg s \vee ps)$ • t

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- SAT modeling

p cnf 5 6
-1 -2 0
2 3 0
-2 -4 -5 0

-2 4 3 0
-3 4 0
5 0
SAT: -1 -2 3 4 5 0

→ Add: 1 2 -3 -4 -5 0
SAT: 1 -2 3 4 5 0

SAT modeling: Propositional logic - 4

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- $ts \implies (ps \vee s) \equiv (\neg ts \vee ps \vee s)$ • $s \implies ps \equiv (\neg s \vee ps)$ • t

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- SAT modeling

p cnf 5 6
-1 -2 0
2 3 0
-2 -4 -5 0

-2 4 3 0
-3 4 0
5 0
SAT: -1 -2 3 4 5 0

Add: 1 2 -3 -4 -5 0
SAT: 1 -2 3 4 5 0
Add: -1 2 -3 -4 -5 0

SAT modeling: Propositional logic - 4

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 \end{aligned}$$

- Goal (\mathcal{G}): All satisfying solutions

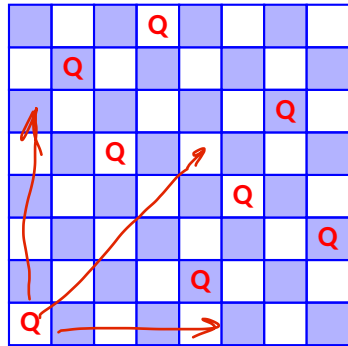
- SAT modeling

p cnf 5 6
 -1 -2 0
 2 3 0
 -2 -4 -5 0

-2 4 3 0
 -3 4 0
 5 0
 SAT: -1 -2 3 4 5 0

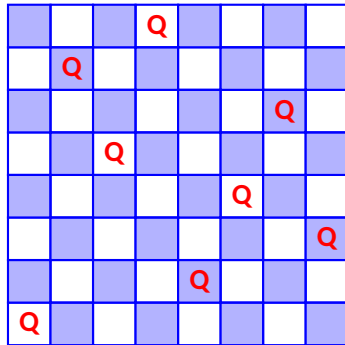
Add: 1 2 -3 -4 -5 0
 SAT: 1 -2 3 4 5 0
 Add: -1 2 -3 -4 -5 0
 UNSAT

SAT modeling: 8-queens



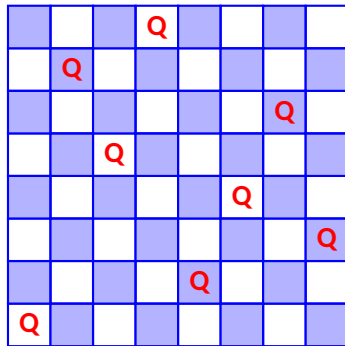
SAT modeling: 8-queens

- Define x_{ij} as (i, j) th cell contains a queen



SAT modeling: 8-queens

- Define x_{ij} as (i, j) th cell contains a queen
- Constraints

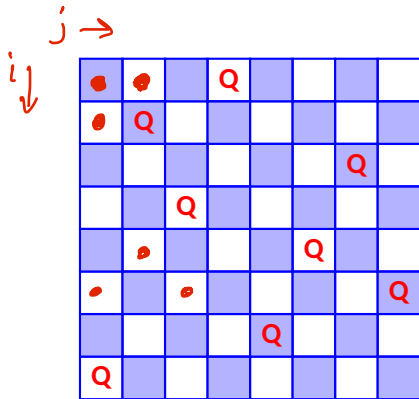


SAT modeling: 8-queens

- Define x_{ij} as (i, j) th cell contains a queen

- Constraints

- $x_{ij} \Rightarrow \neg x_{ij'} \text{ (row)}$
- $x_{ij} \Rightarrow \neg x_{i'j} \text{ (column)}$
- $x_{ij} \Rightarrow \neg x_{(i+k)(j+k)} \text{ (diagonal)}$
- $x_{ij} \Rightarrow \neg x_{(i+k)(j-k)} \text{ (diagonal)}$
- $\bigvee_j x_{ij} \text{ (column)}$
- $\bigvee_i x_{ij} \text{ (row)}$

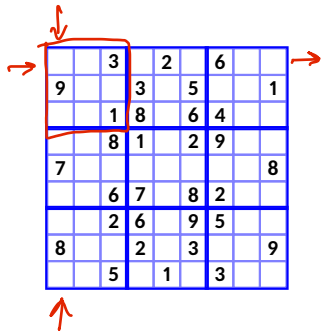


SAT modeling: Sudoku

		3		2		6		
9			3		5			1
		1	8		6	4		
		8	1		2	9		
7								8
		6	7		8	2		
		2	6		9	5		
8			2		3			9
		5		1		3		

SAT modeling: Sudoku

- Define x_{ijk} as (i, j) th cell contains k



SAT modeling: Sudoku

- Define x_{ijk} as (i, j) th cell contains k
- Constraints:

		3		2		6		
9			3		5			1
		1	8		6	4		
		8	1		2	9		
7								8
		6	7		8	2		
		2	6		9	5		
8			2		3			9
		5		1		3		

SAT modeling: Sudoku

- Define x_{ijk} as (i, j) th cell contains k
- Constraints:
 - $x_{ijk} \implies \neg x_{ij'k} \quad \forall i, k, j \neq j' \text{ (same row)}$

•	•	3		2		6		
9			3		5			1
		1	8		6	4		
		8	1		2	9		
7								8
		6	7		8	2		
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8			2		3			9
		5		1		3		

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 - $x_{ijk} \implies \neg x_{i'jk} \quad \forall j, k, i \neq i' \text{ (same column)}$

		3		2		6		
9			3		5			1
		1	8		6	4		
		8	1		2	9		
7								8
		6	7		8	2		
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8			2		3			9
		5		1		3		


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 - $x_{ijk} \implies \neg x_{i'j'k} \quad \forall k, i \neq i', j \neq j', 1 \leq i, i', j, j' \leq 3 \text{ - every block}$

		3		2		6		
9			3		5			1
		1	8		6	4		
		8	1		2	9		
7								8
		6	7		8	2		
		2	6		9	5		
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		5		1		3		

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 - $x_{ijk} \implies \neg x_{i'j'k} \quad \forall k, i \neq i', j \neq j', 1 \leq i, i', j, j' \leq 3$ - every block
 - $\bigvee_j x_{ijk} \quad \forall i, k$ (row)



		3		2		6		
9			3		5			1
		1	8		6	4		
		8	1		2	9		
7								8
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 - $\bigvee_j x_{ijk} \quad \forall i, k$ (row)
 - $\bigvee_i x_{ijk} \quad \forall j, k$ (column)

		3		2		6		
9			3		5			1
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 - $x_{ijk} \implies \neg x_{i'j'k} \quad \forall k, i \neq i', j \neq j', 1 \leq i, i', j, j' \leq 3$ - every block
 - $\bigvee_j x_{ijk} \quad \forall i, k$ (row)
 - $\bigvee_i x_{ijk} \quad \forall j, k$ (column)
 - $\bigvee_{i,j} x_{ijk} \quad \forall k$ every block

		3		2		6		
9			3		5			1
		1	8		6	4		
		8	1		2	9		
7								8
		6	7		8	2		
		2	6		9	5		
8			2		3			9
		5		1		3		

SAT modeling: Sudoku

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 - $\bigvee_j x_{ijk} \quad \forall i, k$ (row)
 - $\bigvee_i x_{ijk} \quad \forall j, k$ (column)
 - $\bigvee_{i,j} x_{ijk} \quad \forall k$ every block
 - $x_{ijk} \implies \neg x_{ijk'} \quad \forall i, j, k \neq k'$ (same cell)



		3		2		6		
9			3		5			1
		1	8		6	4		
		8	1		2	9		
7								8
		6	7		8	2		
		2	6		9	5		
8			2		3			9
		5		1		3		

x_{111} x_{112}

SAT modeling: Sudoku

- Define x_{ijk} as (i, j) th cell contains k
- Constraints:
 - $x_{ijk} \implies \neg x_{ij'k} \quad \forall i, k, j \neq j'$ (same row)
 - $x_{ijk} \implies \neg x_{i'jk} \quad \forall j, k, i \neq i'$ (same column)
 - $x_{ijk} \implies \neg x_{i'j'k} \quad \forall k, i \neq i', j \neq j', 1 \leq i, i', j, j' \leq 3$ - every block
 - $\bigvee_j x_{ijk} \quad \forall i, k$ (row)
 - $\bigvee_i x_{ijk} \quad \forall j, k$ (column)
 - $\bigvee_{i,j} x_{ijk} \quad \forall k$ every block
 - $x_{ijk} \implies \neg x_{ijk'} \quad \forall i, j, k \neq k'$ (same cell)
 - $\bigvee_k x_{ijk} \quad \forall i, j$ (every cell)

		3		2		6		
9			3		5			1
		1	8		6	4		
		8	1		2	9		
7								8
		6	7		8	2		
		2	6		9	5		
8			2		3			9
		5		1		3		

SAT modeling: Sudoku

- Define x_{ijk} as (i, j) th cell contains k
- Constraints:
 - $x_{ijk} \implies \neg x_{ij'k} \quad \forall i, k, j \neq j'$ (same row)
 - $x_{ijk} \implies \neg x_{i'jk} \quad \forall j, k, i \neq i'$ (same column)
 - $x_{ijk} \implies \neg x_{i'j'k} \quad \forall k, i \neq i', j \neq j', 1 \leq i, i', j, j' \leq 3$ - every block
 - $\bigvee_j x_{ijk} \quad \forall i, k$ (row)
 - $\bigvee_i x_{ijk} \quad \forall j, k$ (column)
 - $\bigvee_{i,j} x_{ijk} \quad \forall k$ every block
 - $x_{ijk} \implies \neg x_{ijk'} \quad \forall i, j, k \neq k'$ (same cell)
 - $\bigvee_k x_{ijk} \quad \forall i, j$ (every cell)
 - x_{133}, x_{176}, \dots

		3		2		6		
9			3		5			1
		1	8		6	4		
		8	1		2	9		
7								8
		6	7		8	2		
		2	6		9	5		
8			2		3			9
		5		1		3		

SAT modeling: Langford sequence

- Given the bag of numbers $\{1, 1, 2, 2, 3, 3, \dots, n, n\}$, can they be arranged in a sequence $L(n)$ such that for $1 \leq i \leq n$ there are i numbers between the two occurrences of i ?
 - $L(4) = 41312432$
 - $L(3) = ?$

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	1	2	3	4	5	6
x_1	1		1			
x_2		1		1		
x_3			1		1	
x_4				1		1
x_5	2			2		
x_6		2			2	
x_7			2			2
x_8	3				3	
x_9		3				3

SAT modeling: Langford sequence

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- $L(4) = 41312432$

- $L(3) = ?$

- Constraints:

- $x_1 \vee x_2 \vee x_3 \vee x_4$

- $x_k \implies \neg x_{k'} \quad 1 \leq k < k' \leq 4$

- Similarly for the other numbers

- $x_1 \vee x_5 \vee x_8$ ✓

- $x_1 \implies \neg x_5, \dots$

- Similarly for the other columns

	1	2	3	4	5	6
x_1	<u>1</u>		1			
x_2		1		1		
x_3			1		1	
x_4				1		1
x_5		<u>2</u>		2		
x_6		2			2	
x_7			2			2
x_8	<u>3</u>				3	
x_9		3				3

Thank you!